Worksheet: Arithmetic Expressions – Some Answers

(1) Using the rules from the lectures:

$$\frac{\frac{4 \downarrow 4}{(B-NUM)} \frac{1 \downarrow 1}{1 \downarrow 1}}{(4+1) \downarrow 5}_{(B-ADD)} \frac{\frac{2 \downarrow 2}{(B-NUM)} \frac{2 \downarrow 2}{2 \downarrow 2}}{(2+2) \downarrow 4}_{(B-ADD)}}{(2+2) \downarrow 4}_{(B-ADD)}$$

(2) To handle multiplication, we simply add a single rule to our existing system. In addition to the axiom and the rule for + we have the rule

(B-MULT)
$$\frac{E_1 \Downarrow \mathbf{n}_1 \quad E_2 \Downarrow \mathbf{n}_2}{(E_1 \times E_2) \Downarrow \mathbf{n}_3} n_3 = \mathsf{mult}(n_1, n_2)$$

(3) The proof consists of four uses of the axiom, two uses of the rule for + and one use of the new rule for ×:

$$\frac{\overline{3 \downarrow 3}}{3 + 2 \downarrow 5}^{(B-NUM)} \underbrace{\overline{2 \downarrow 2}^{(B-NUM)}}_{(B-ADD)} \underbrace{\overline{1 \downarrow 1}^{(B-NUM)} \overline{4 \downarrow 4}^{(B-NUM)}}_{1 + 4 \downarrow 5}_{(B-ADD)}$$

$$\underbrace{((3 + 2) \times (1 + 4)) \downarrow 25}^{(B-NUM)}$$

(4) The obvious rule to add for subtraction is

$$_{\text{(B-SUB)}} \frac{E_1 \Downarrow \mathbf{n}_1 \quad E_2 \Downarrow \mathbf{n}_2}{(E_1 - E_2) \Downarrow \mathbf{n}_3} n_3 = \min(n_1, n_2)$$

However, in general this won't work, because if n_1 is 3 and n_2 is 7, then $n_3 = -4$, and we do not have a corresponding numeral n_3 .

One solution is to say that subtraction "gets stuck" when a negative value is needed. You can do this by making no rule available in the nasty case; this is done by adding a side condition.

(B-SUB)
$$\frac{E_1 \Downarrow n_1 \quad E_2 \Downarrow n_2}{(E_1 - E_2) \Downarrow n_3} n_3 = \min(n_1, n_2) \text{ and } n_1 \ge n_2$$

With this approach, there is no numeral **n** for which $(3 - 7) \Downarrow n$, so this expression has *no* final answer at all.

Other solutions are possible: one could extend the semantics to allow expressions to signal errors, and then a "bad" subtraction would signal an error. If you choose this route, every rule of the semantics needs to be reconsidered in case errors might make a difference.

You might also decide that such subtractions default to returning 0, but then you have of beware of strange things like

$$(3-7)+4 \Downarrow 4$$

Answers

while

 $(3+4) - 7 \Downarrow 0.$

This is probably not a good idea.

(5) One possible *E* is 3 + (4 + 3). The axiom (s-ADD) for the small-step semantics gives us

$$1 + 2 \rightarrow_{lr} 3$$

and the rule $_{(\text{S-LEFT})}$ lets us use this on the left hand side of the expression in question, so the full derivation is

$$\frac{\overline{1+2 \rightarrow_{lr} 3}^{(s-add)}}{((1+2)+(4+3)) \rightarrow_{lr} (3+(4+3))}^{(s-add)}$$

Question: Are there any other expression *E* different from 3 + (4 + 3) for which a left-to-right derivation can be found ?

(6) The full evaluation sequence is

$$\begin{array}{rcl} ((1+2)+(4+3)) & \to_{\mathrm{lr}} & (3+(4+3)) \\ & \to_{\mathrm{lr}} & (3+7) \\ & \to_{\mathrm{lr}} & \mathbf{10}. \end{array}$$

We have already seen the derivation of the first step. The axiom for the small step semantics (s-ADD) allows us to derive

$$4 + 3 \rightarrow_{lr} 7$$

Since 3 is a numeral this can be used in an application of the second rule, (s-N.RIGHT) to give the following derivation of the second step:

$$\frac{\overline{4+3\rightarrow_{lr}7}}{3+(4+3)\rightarrow_{lr}(3+7)}^{(s-add)}$$
(s-n.right)

The derivation of the final step is simpler. It is an application of the axiom:

$$\overline{3+7 \rightarrow_{lr} 10}^{(s-add)}$$

(7) Every derivation in the left-to-right semantics is also a derivation in the standard semantics. So from Question (5) we know

$$((1+2)+(4+3)) \rightarrow 3+(4+3)$$

But using the more general rule (S-RIGHT) we can also derive

$$((1+2)+(4+3)) \rightarrow (1+2)+7$$

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Here is the derivation:

$$\frac{\overline{4+3 \to 7}^{(s-add)}}{((1+2)+(4+3)) \to (1+2)+7}$$
(s-right)

It turns out that these are the only two possible *E*, namely (1+2)+7 and 3+(4+3).

(8) The left-to-right small step semantics uses only three rules, (s-LEFT), (s-N.RIGHT) and (s-ADD). Therefore if an expression is to evaluate to ((1 + 2) + 7) in one step, the step must be derived from an instance of the first rule, (s-LEFT), since the axiom (s-ADD) only lets us derive steps that end in a numeral, and the second rule (s-N.RIGHT) only lets us derive steps that lead to expressions of the form n + E.

Therefore we know that any such expression must have the form $(E_1 + 7)$, where $E_1 \rightarrow_{lr} (1 + 2)$.

So we need to find such an E_1 . Using the first rule (S-LEFT) we can derive

$$((0+1)+2) \rightarrow (1+2)$$

So one possible *E* is (0 + 1) + 2) + 7. The full derivation consists of one use of the axiom followed by two applications of the first rule:

$$\frac{\overline{(\mathbf{0}+1) \rightarrow 1}^{(\text{s-add})}}{\overline{(\mathbf{0}+1)+2} \rightarrow_{\text{lr}} 1+2}^{(\text{s-left})}}{\overline{((\mathbf{0}+1)+2)+7} \rightarrow_{\text{lr}} (1+2)+7}^{(\text{s-left})}$$

(9) From the reasoning in the previous answer we know that they all have the form $(E_1 + 7)$, where $E_1 \rightarrow (1 + 2)$. So we need to know all the possible E_1 such that $E_1 \rightarrow (1 + 2)$.

This step can be derived using either the first rule (s-LEFT) or the second rule (s-RIGHT). Here are all the suitable expressions E_1 :

- ((0 + 1) + 2), using (s-left)
- ((1 + 0) + 2), using (s-left)
- (1 + (0 + 2)), using (s-N.RIGHT)
- (1 + (1 + 1)), using (s-N.RIGHT)
- (1 + (2 + 0)), using (s-N.RIGHT)

So the only expressions that can evaluate to ((1 + 2) + 7) in one step are

- (((0+1)+2)+7)
- (((1+0)+2)+7)
- ((1 + (0 + 2)) + 7)
- ((1 + (1 + 1)) + 7)

Answers

• ((1 + (2 + 0)) + 7).

(10) If $E \rightarrow ((1+2)+7)$ but $E \rightarrow_{lr} ((1+2)+7)$ then the more general rule (s-RIGHT) must be used in the derivation of the first judgement. So one possible *E* is ((1+2)+(2+5)). Here is a derivation:

$$\frac{\overline{2+5 \to 7}^{\text{(s-add)}}}{((1+2)+(2+5)) \to (1+2)+7}^{\text{(s-right)}}$$

(11) ((1+1)+1) takes two steps to get to the final answer 3. The full evaluation sequence is

 $((1+1)+1) \rightarrow_{lr} (2+1) \rightarrow_{lr} 3.$

If we add another 1 to get (((1 + 1) + 1) + 1) then three steps are needed.

In general, the number of steps needed is the same as the number of + symbols since each step reduces the number of + symbols in an expression by one. You will be asked to prove this fact in a later exercise class.