

THE PER MODEL OF ABSTRACT NON-INTERFERENCE

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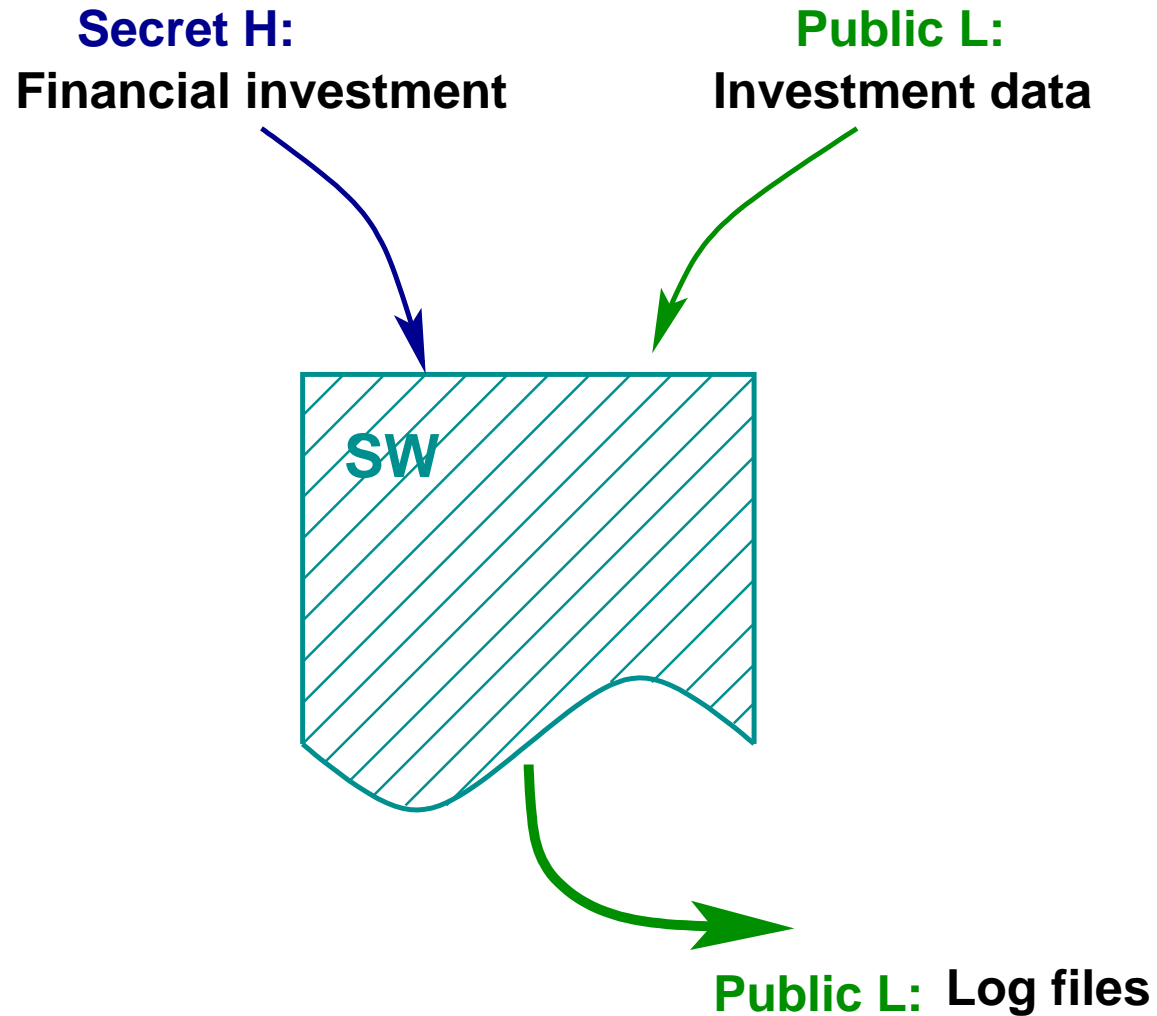
and

Department of Computing and Informatic Sciences

Kansas State University, Manhattan, KS, USA

SAS, September 9th, 2005

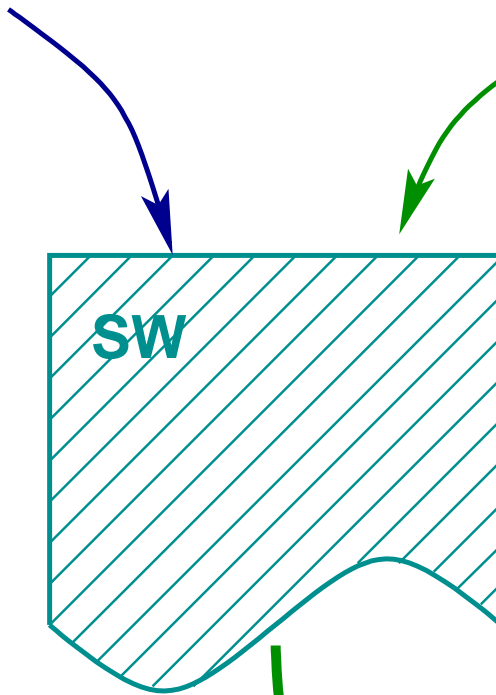
The Problem: Non-Interference



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Secret H:
Financial investment

Public L:
Investment data



Is it secure?

Public L: Log files

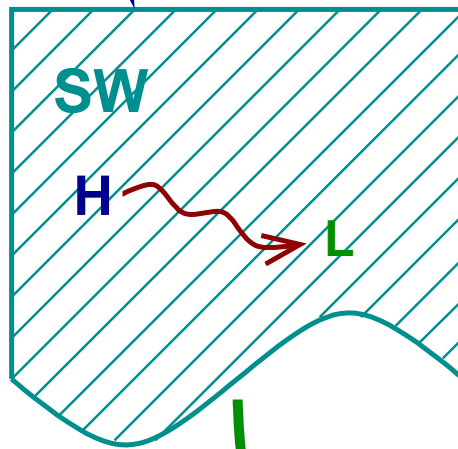


External observer

The Problem: Non-Interference

Secret H:
Financial investment

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Investment data



Is it secure? **NO**

Secret H
Public L: Log files



External observer

Background: The PER model

SECURITY PROPERTY: States which classes have not to interfere with other classes of objects.

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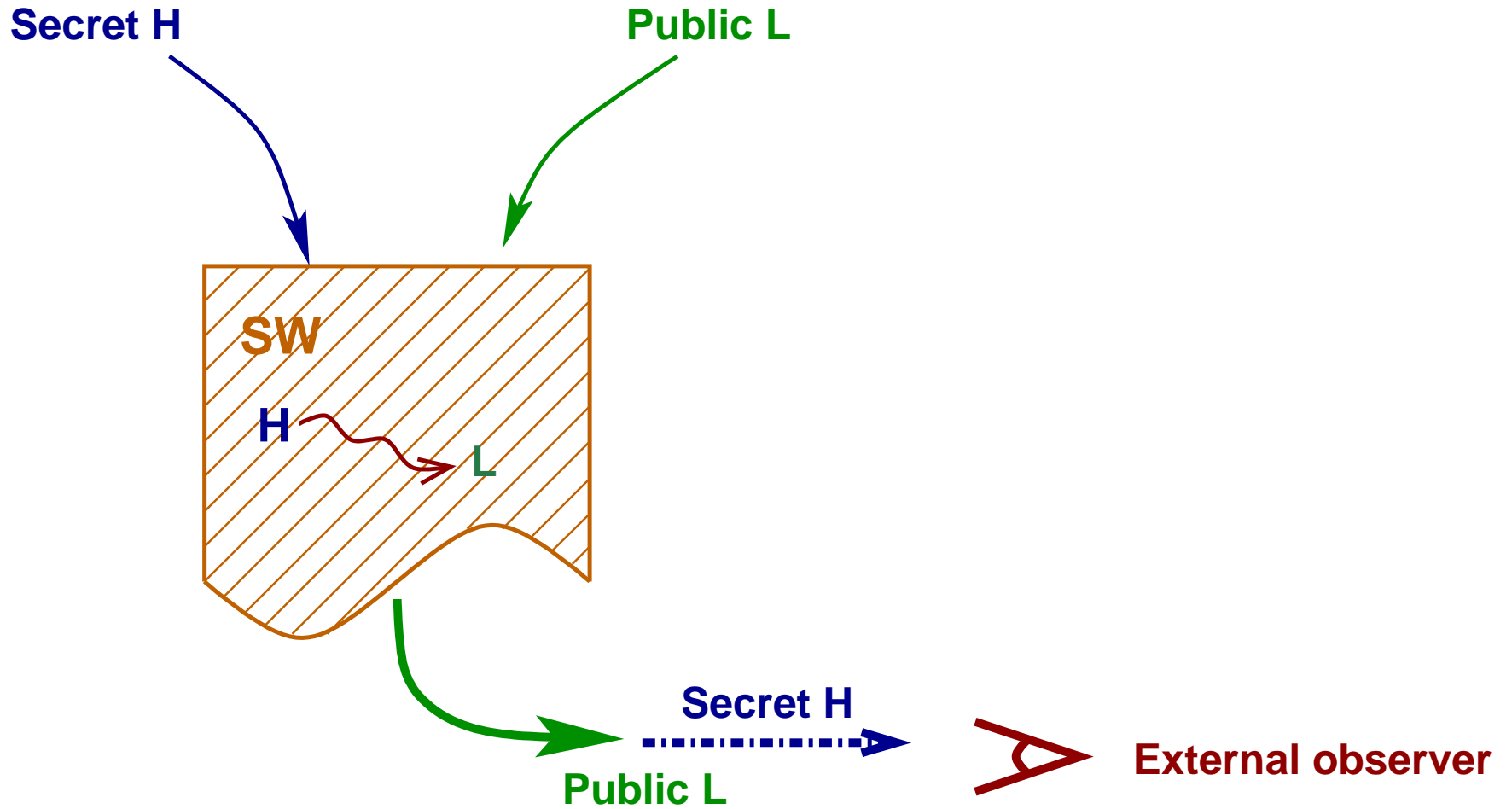
Confinement problem [Lampson'73]: Preventing the results of computations leaking even partial information about the confidential inputs.

⑥ **PER model [Sabelfeld & Sands'01]:**

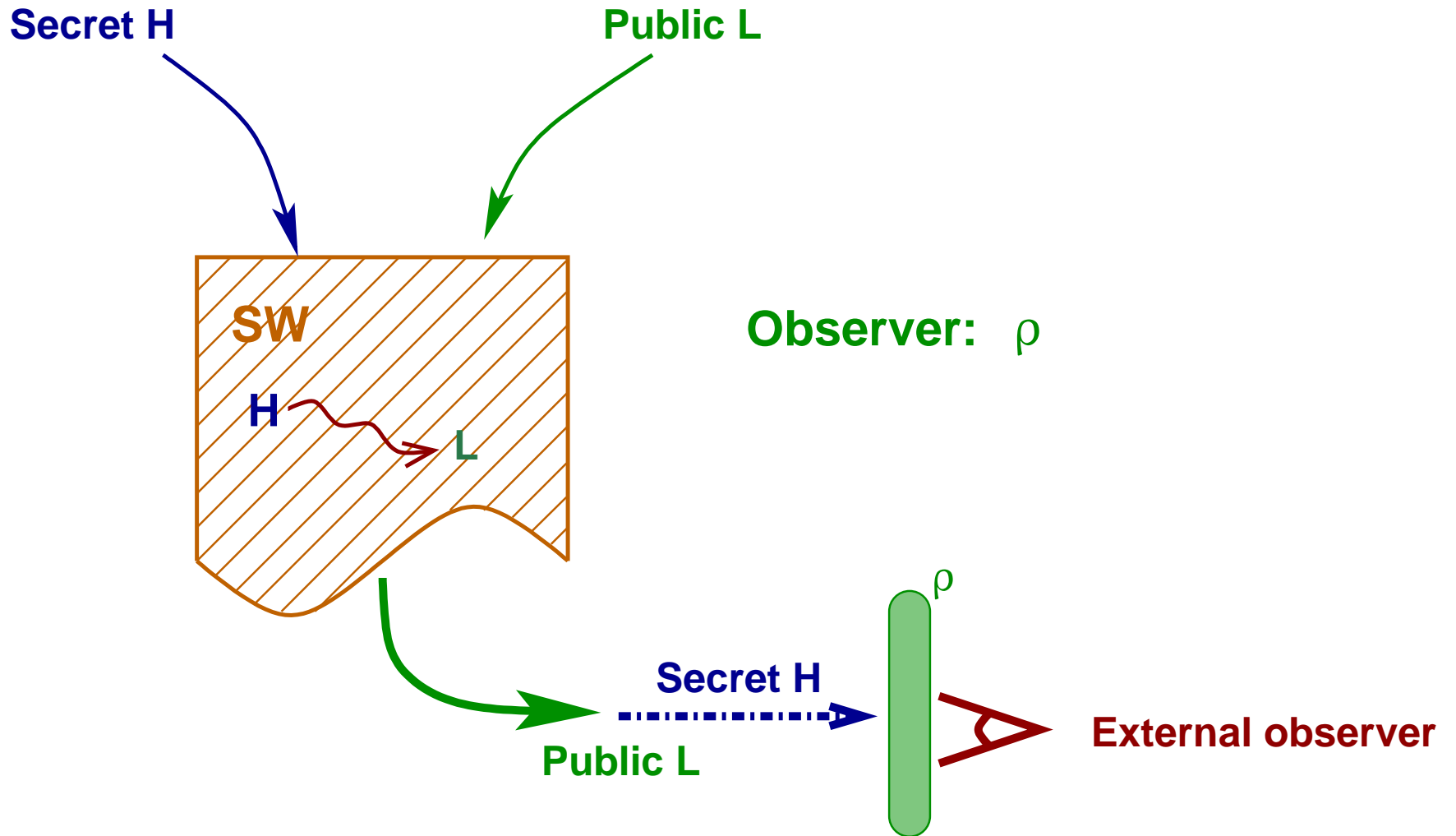
Let $\forall x, x'. x \text{ All } x' \text{ and } x \text{ Id } x' \Leftrightarrow x = x'$.

P is *secure* iff
 $\forall s, t. \langle s^H, s^L \rangle \text{ All } \times \text{ Id } \langle t^H, t^L \rangle \Rightarrow \llbracket P \rrbracket(s) \text{ All } \times \text{ Id } \llbracket P \rrbracket(t)$

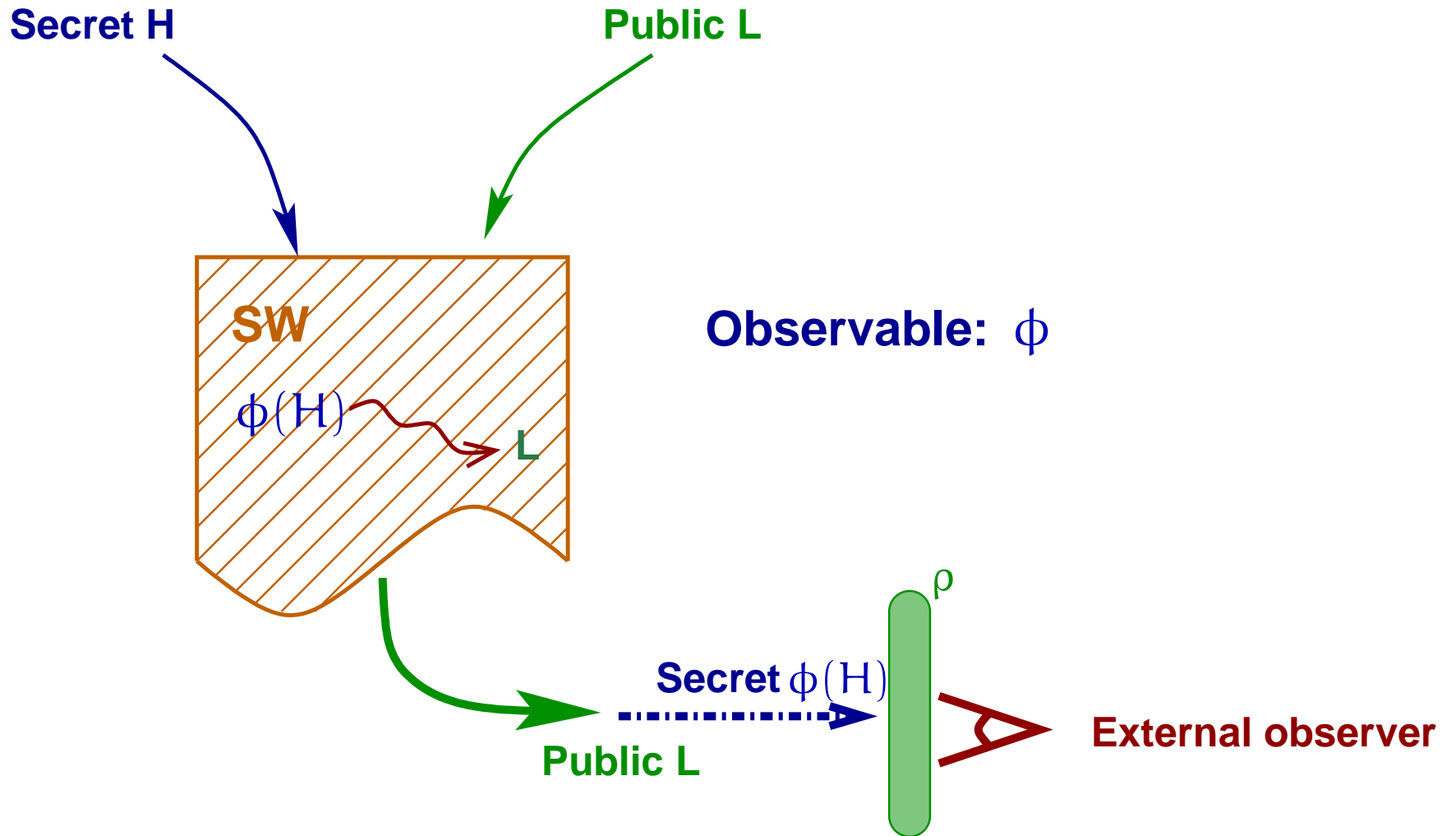
Abstracting Non-Interference



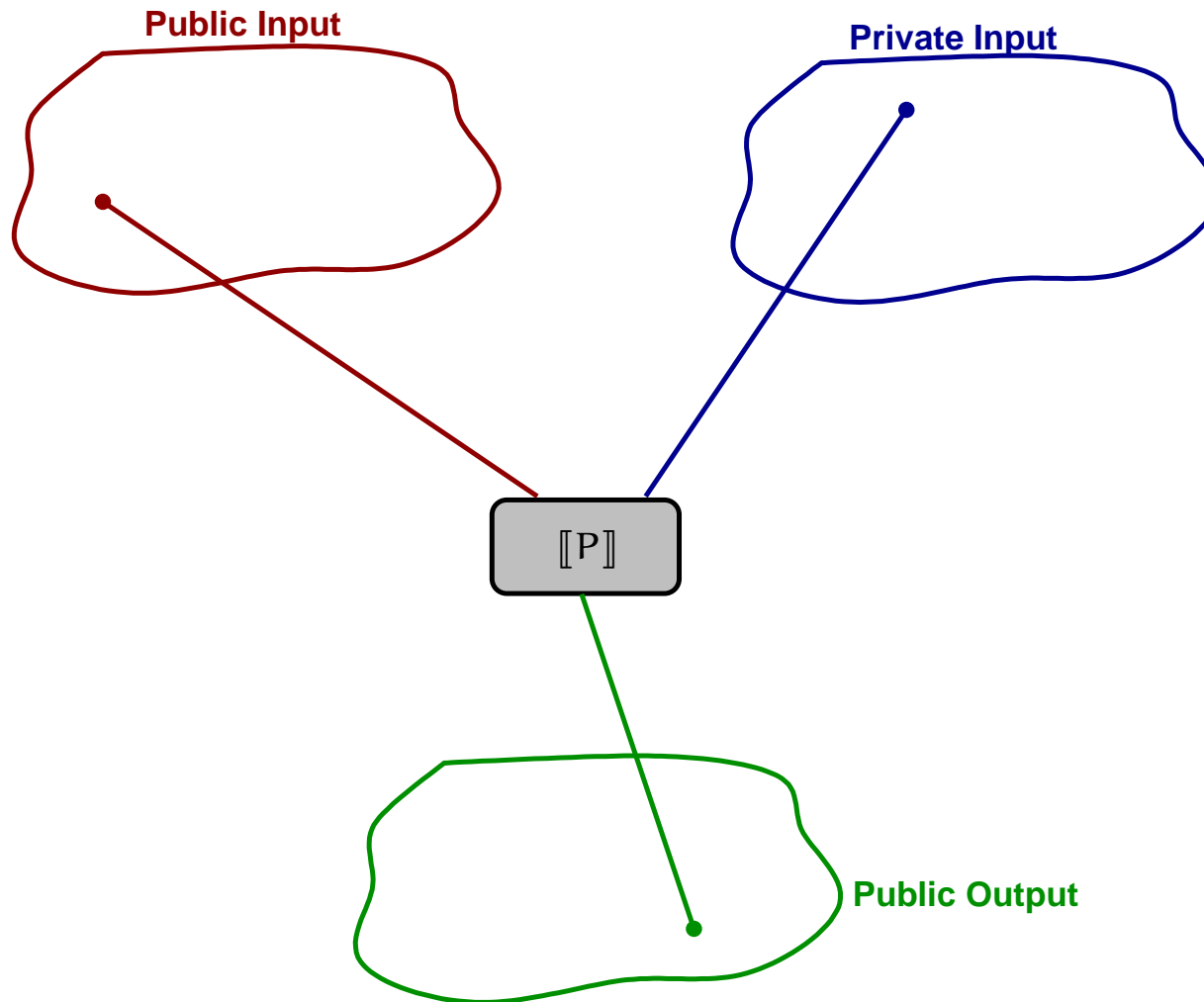
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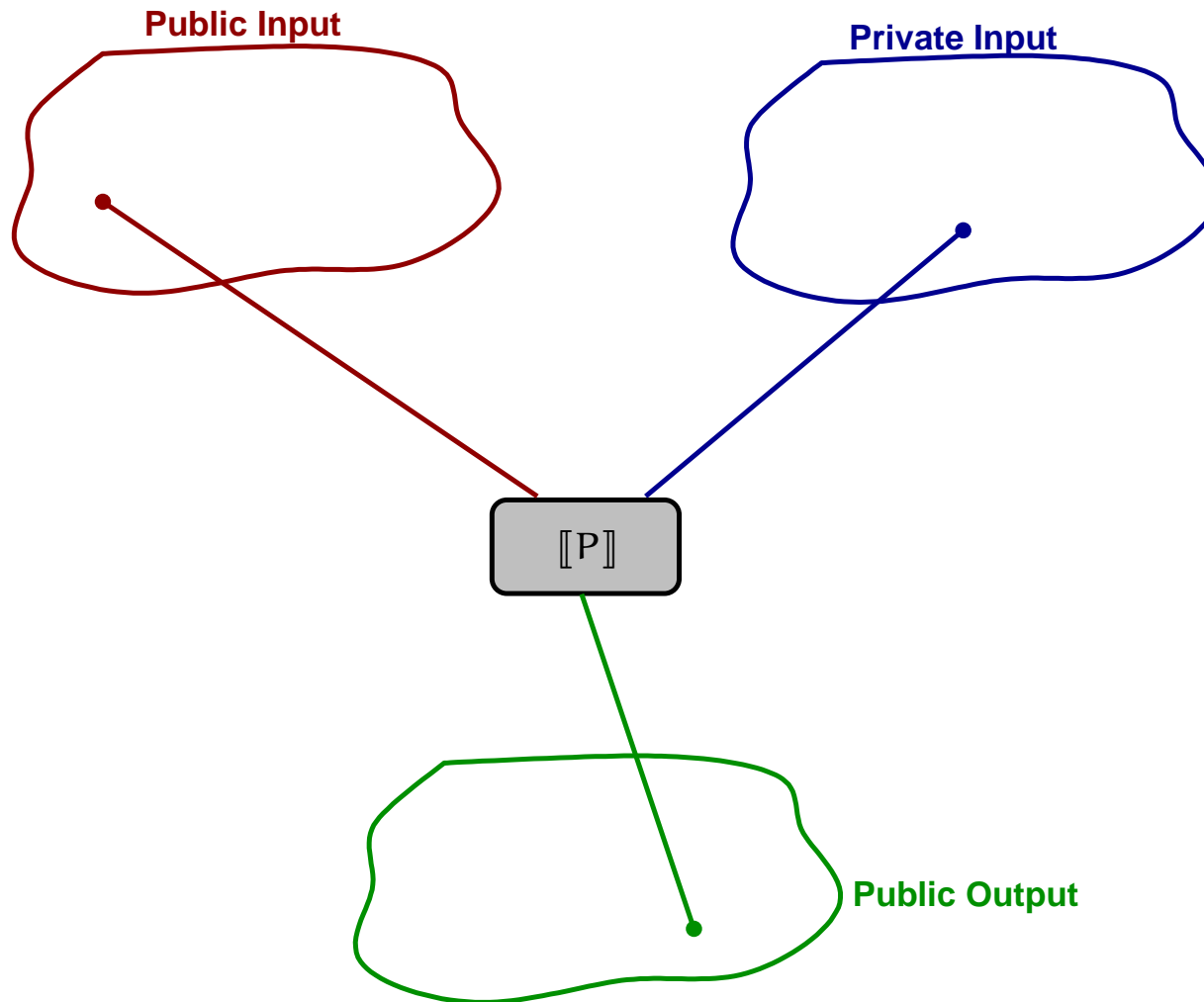


Standard non-interference



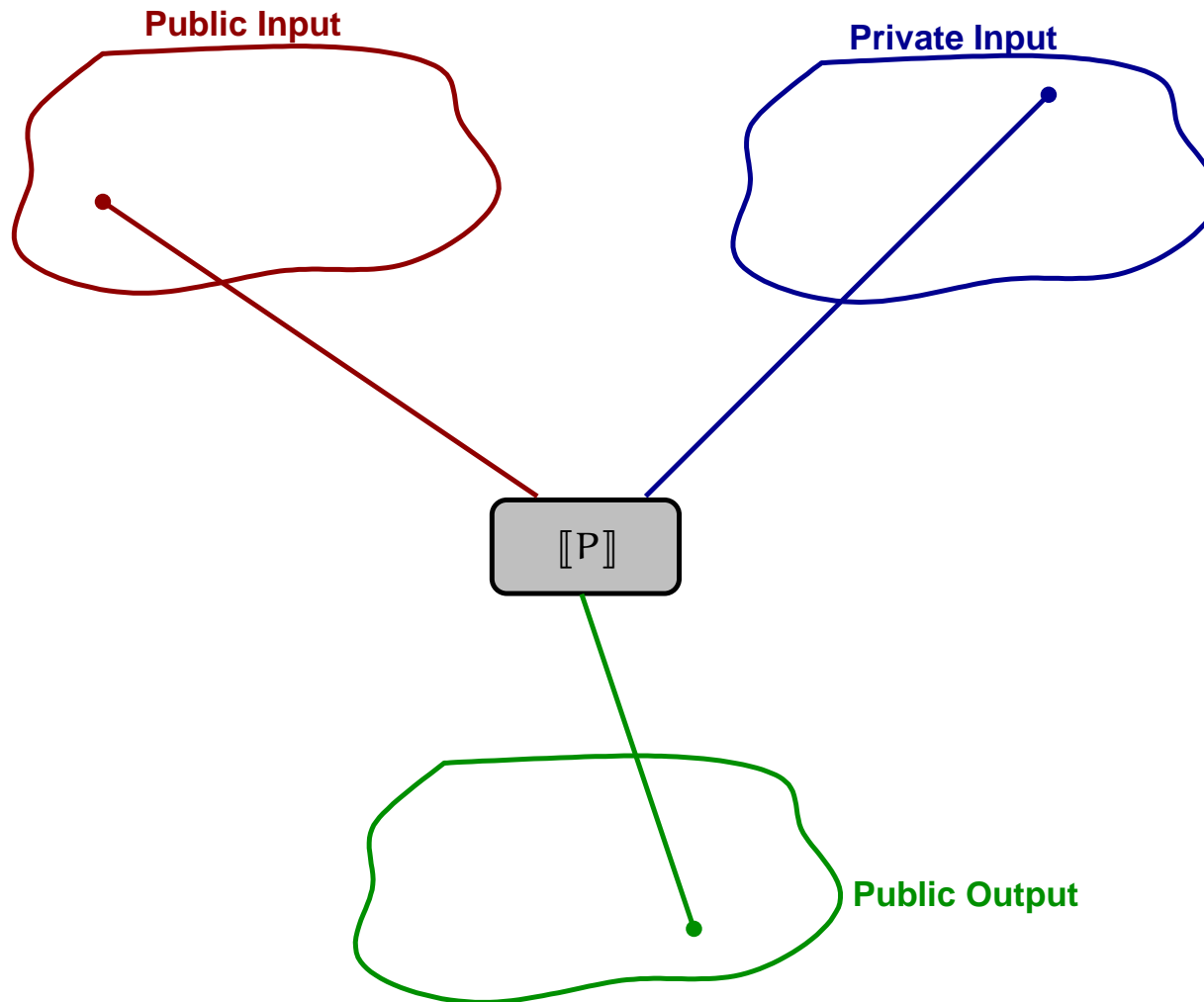
$$\forall l : L, \forall h_1, h_2 : H. [[P]](h_1, l)^L = [[P]](h_2, l)^L$$

Standard non-interference



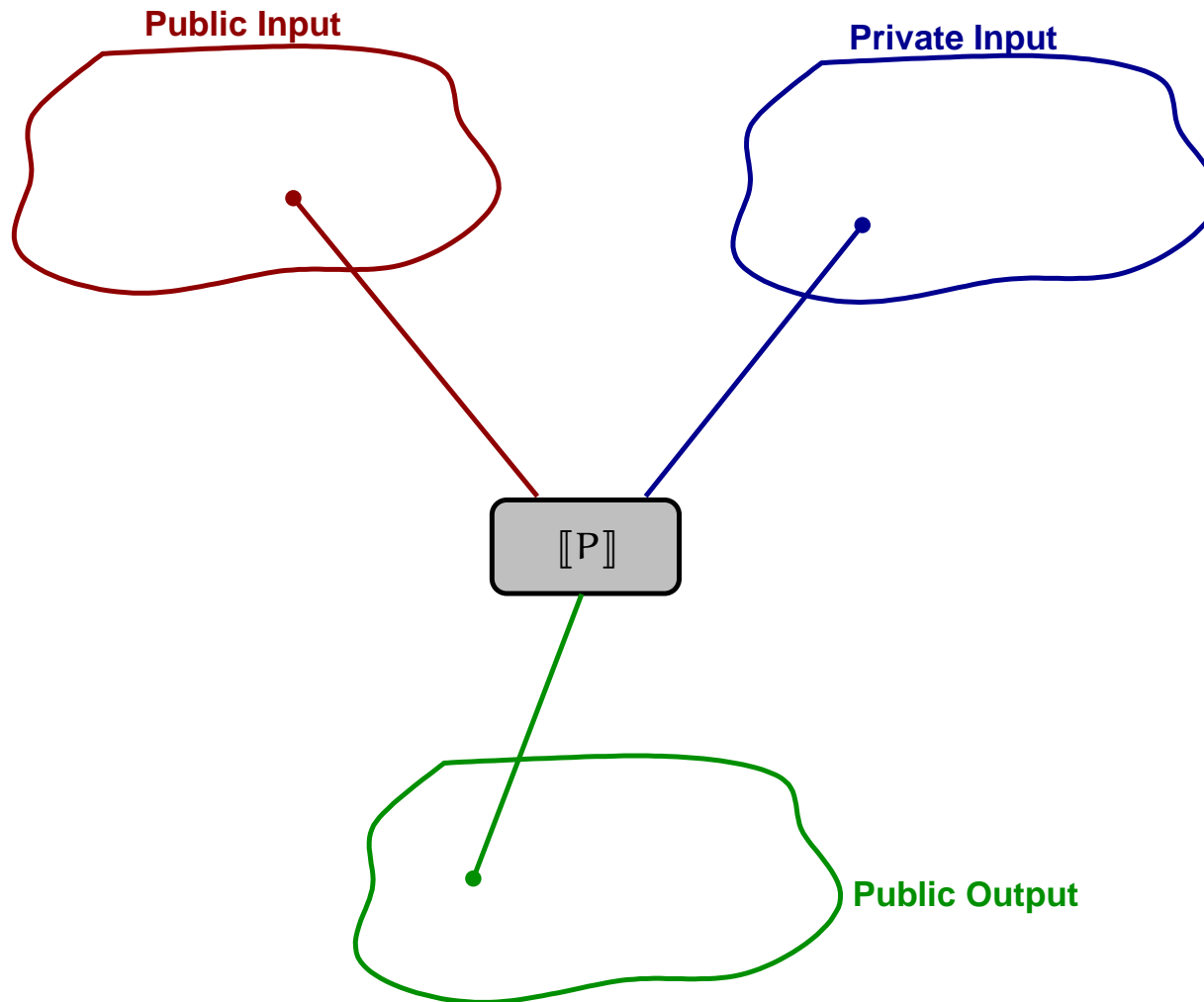
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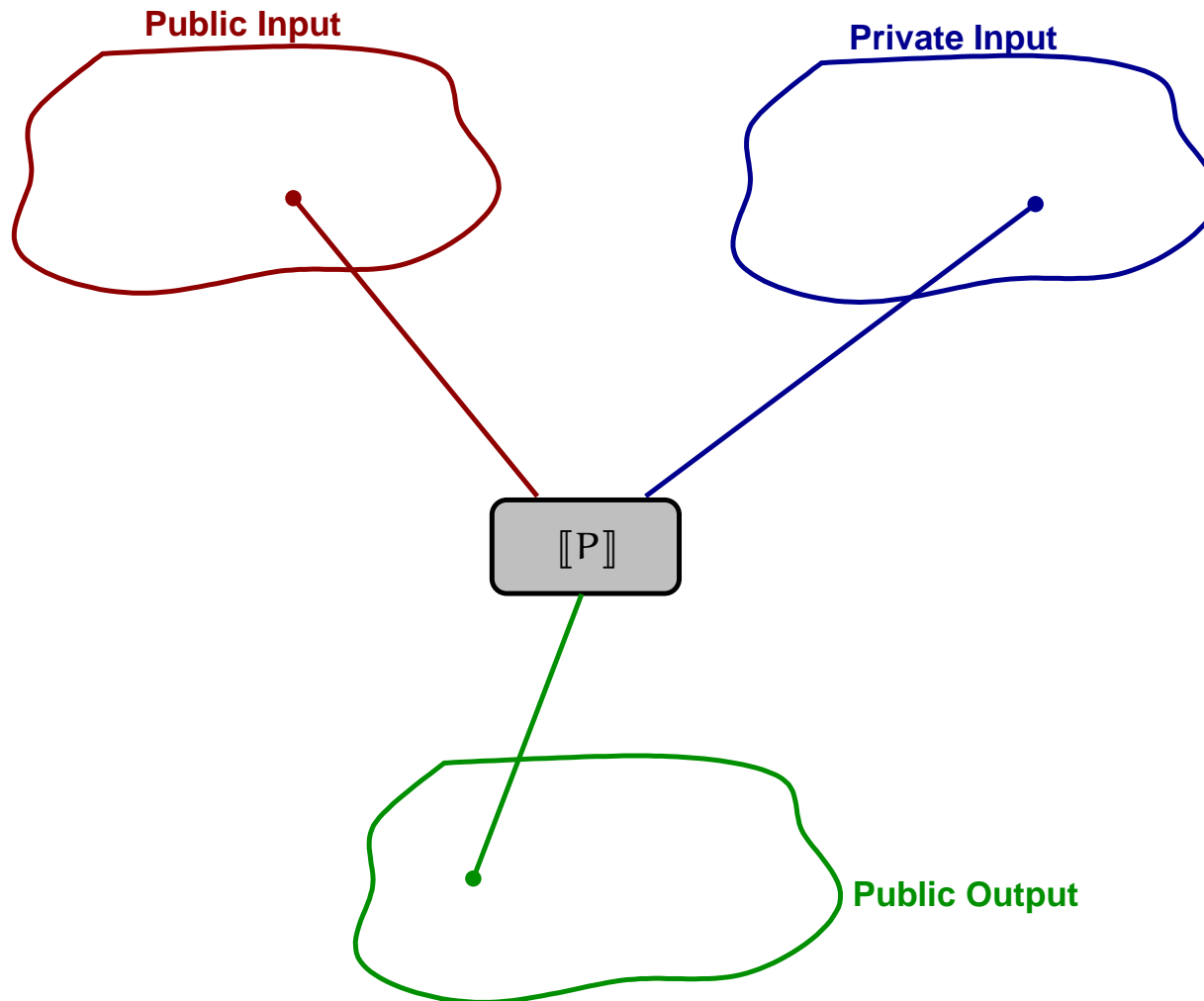
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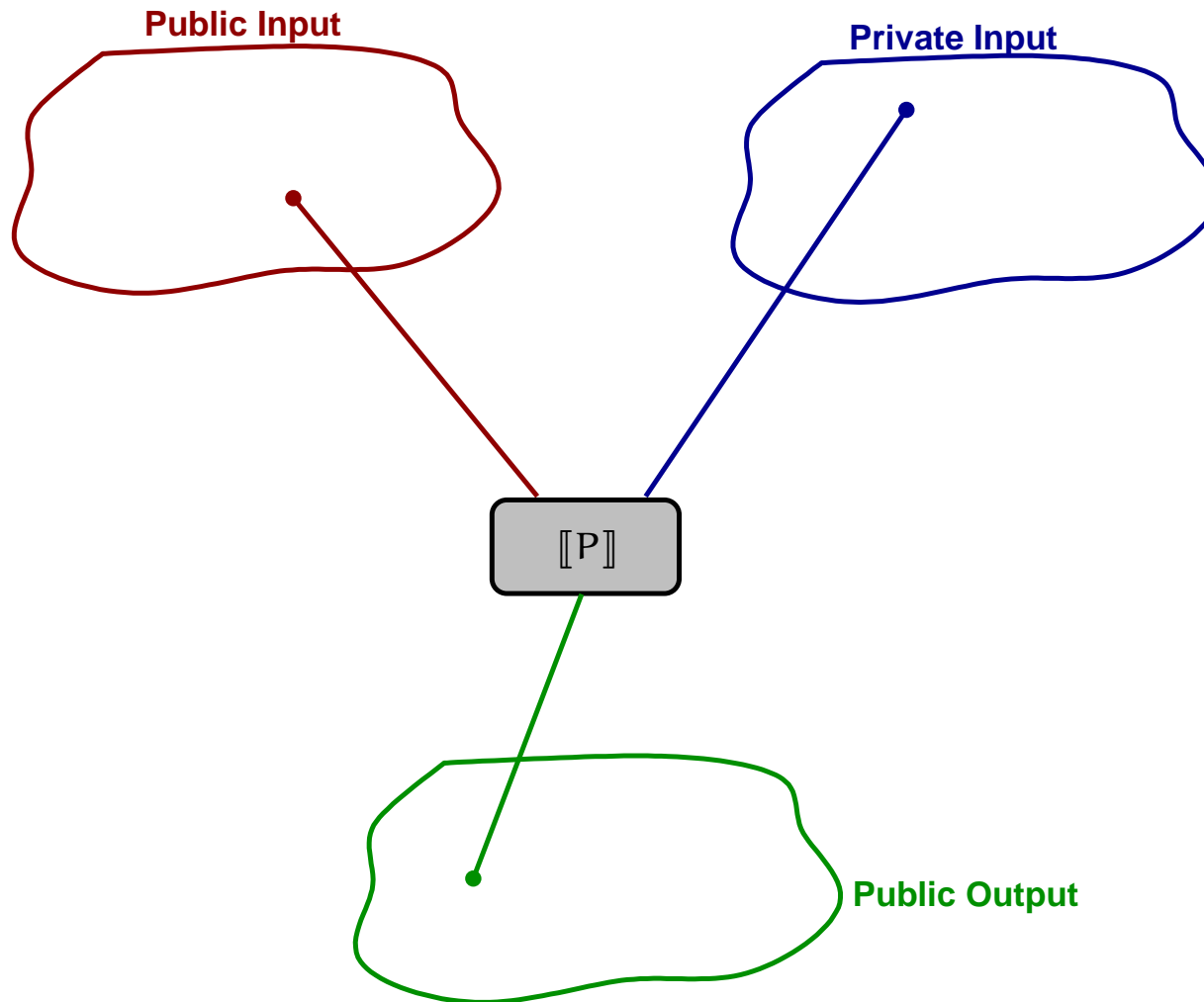
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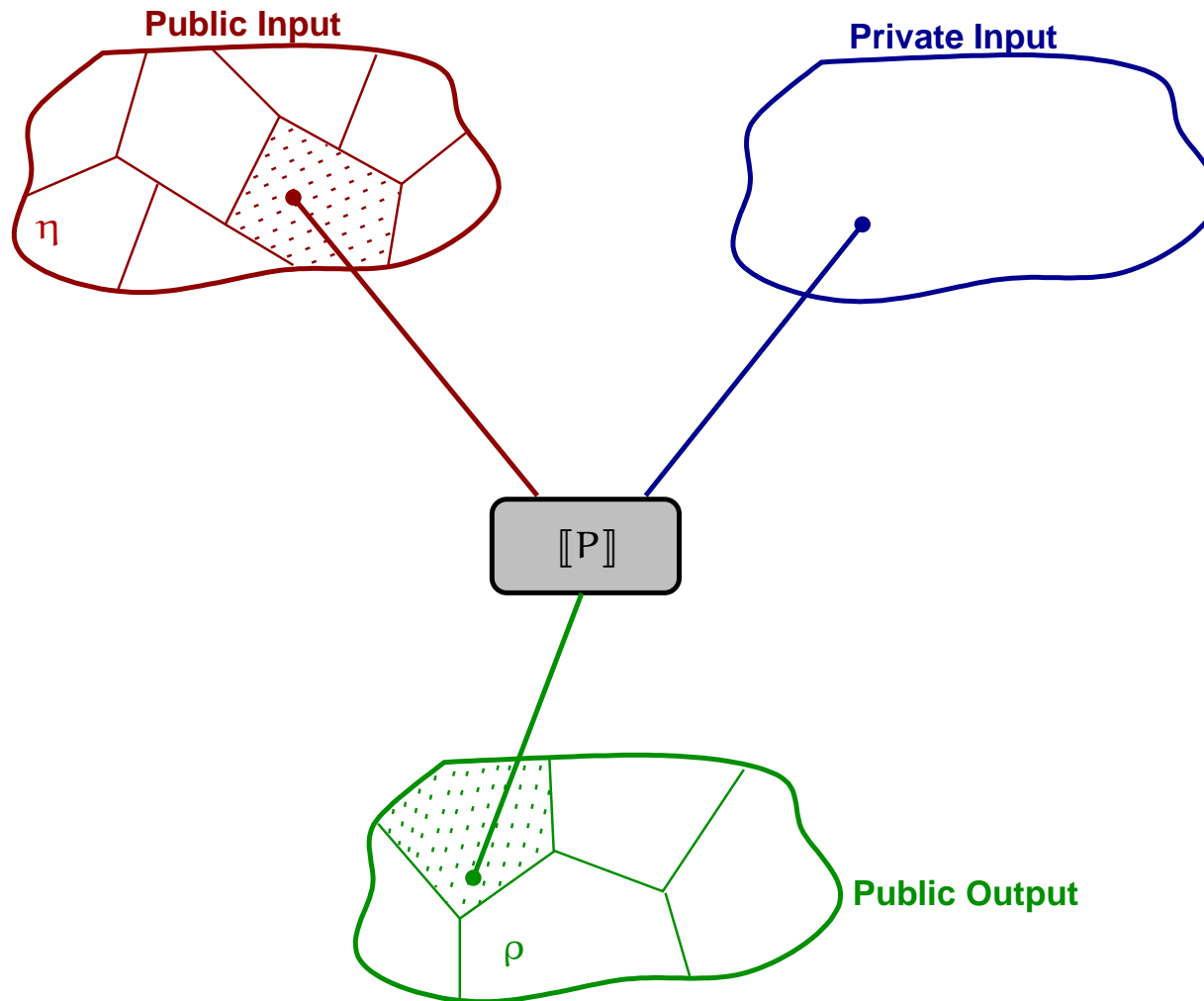
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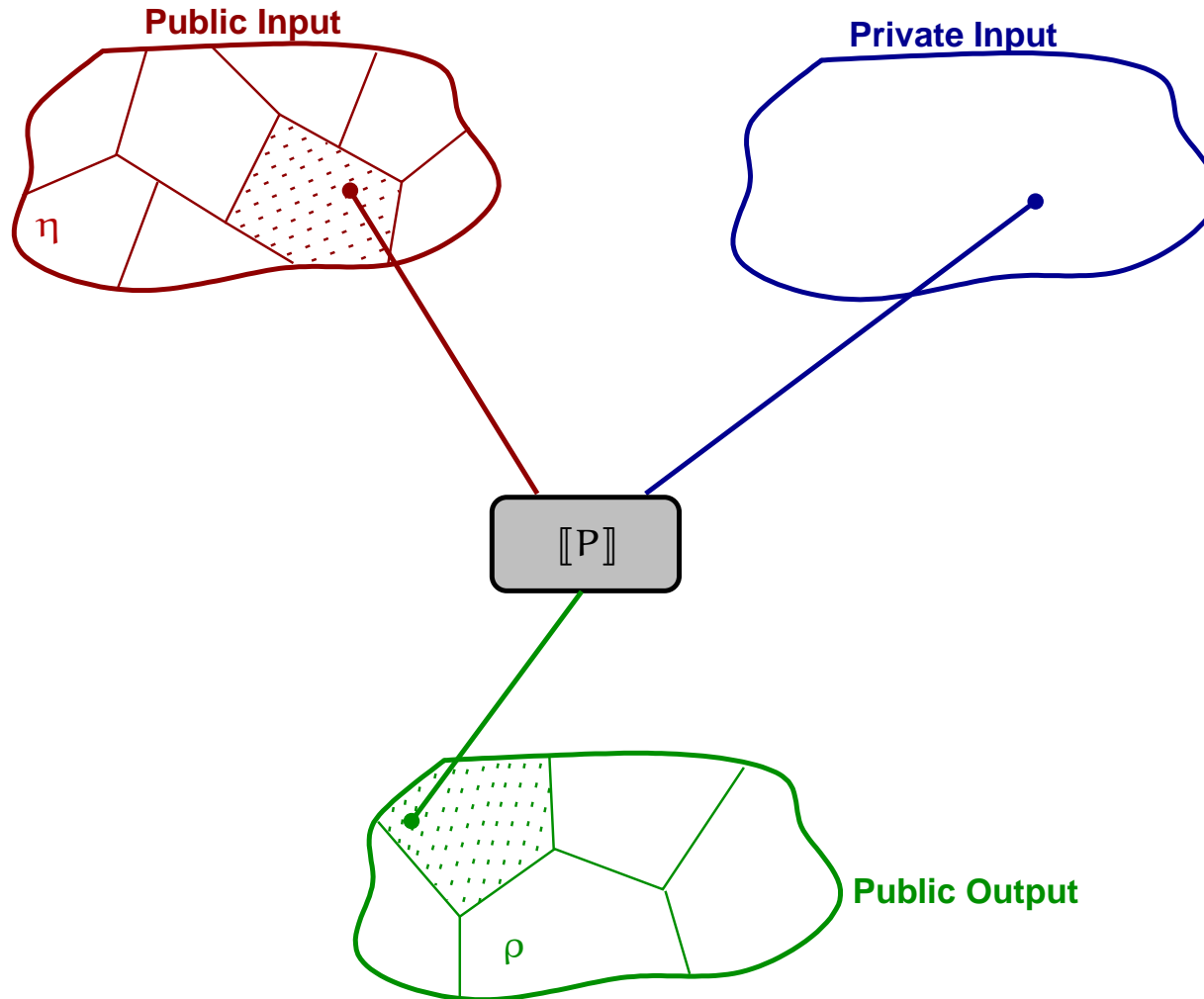
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Abstracting non-interference I: Narrow ANI



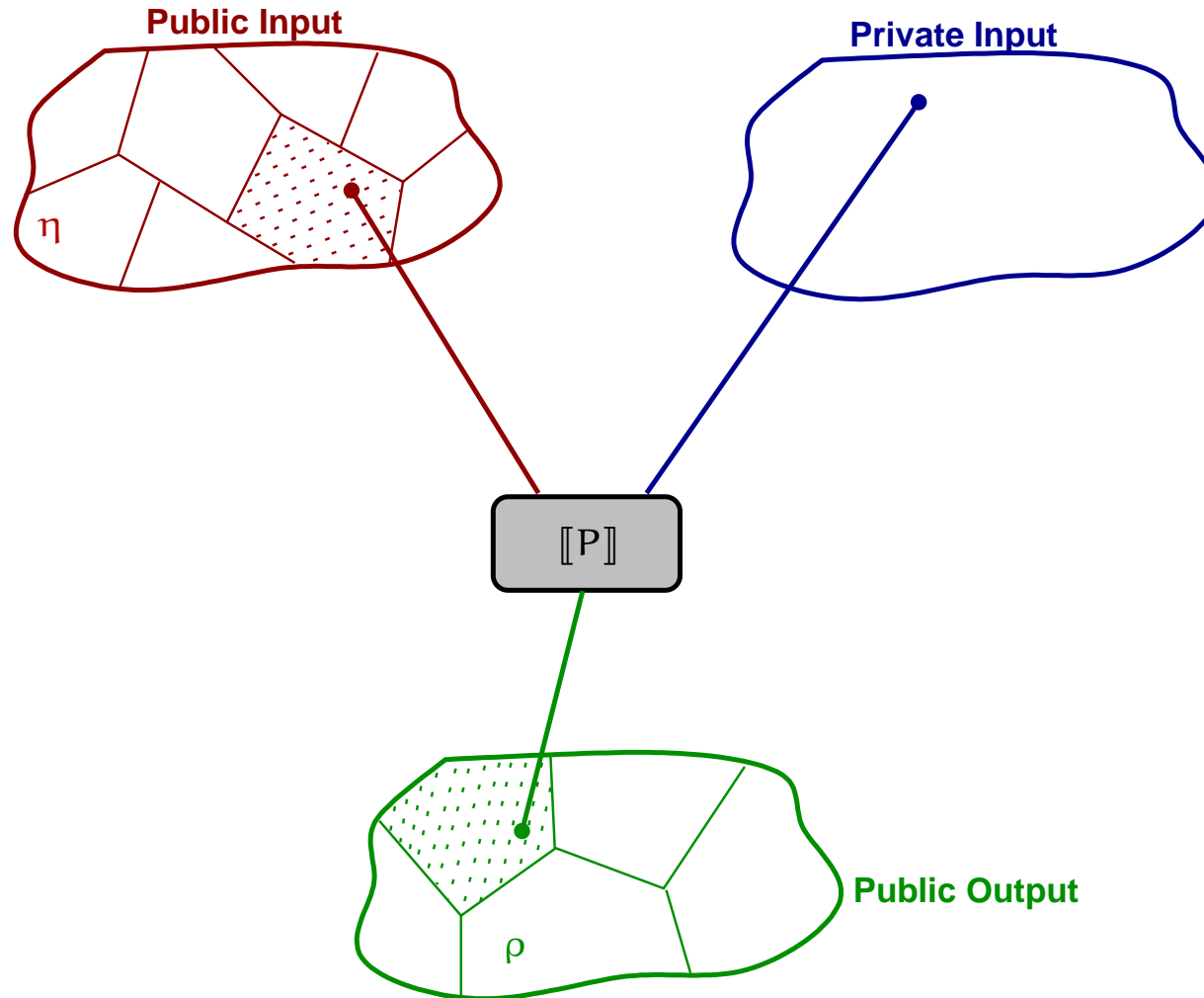
$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)): [\eta]P(\rho): \eta(l_1) = \eta(l_2) \Rightarrow \rho([\![P]\!](h_1, l_1)^L) = \rho([\![P]\!](h_2, l_2)^L)$$

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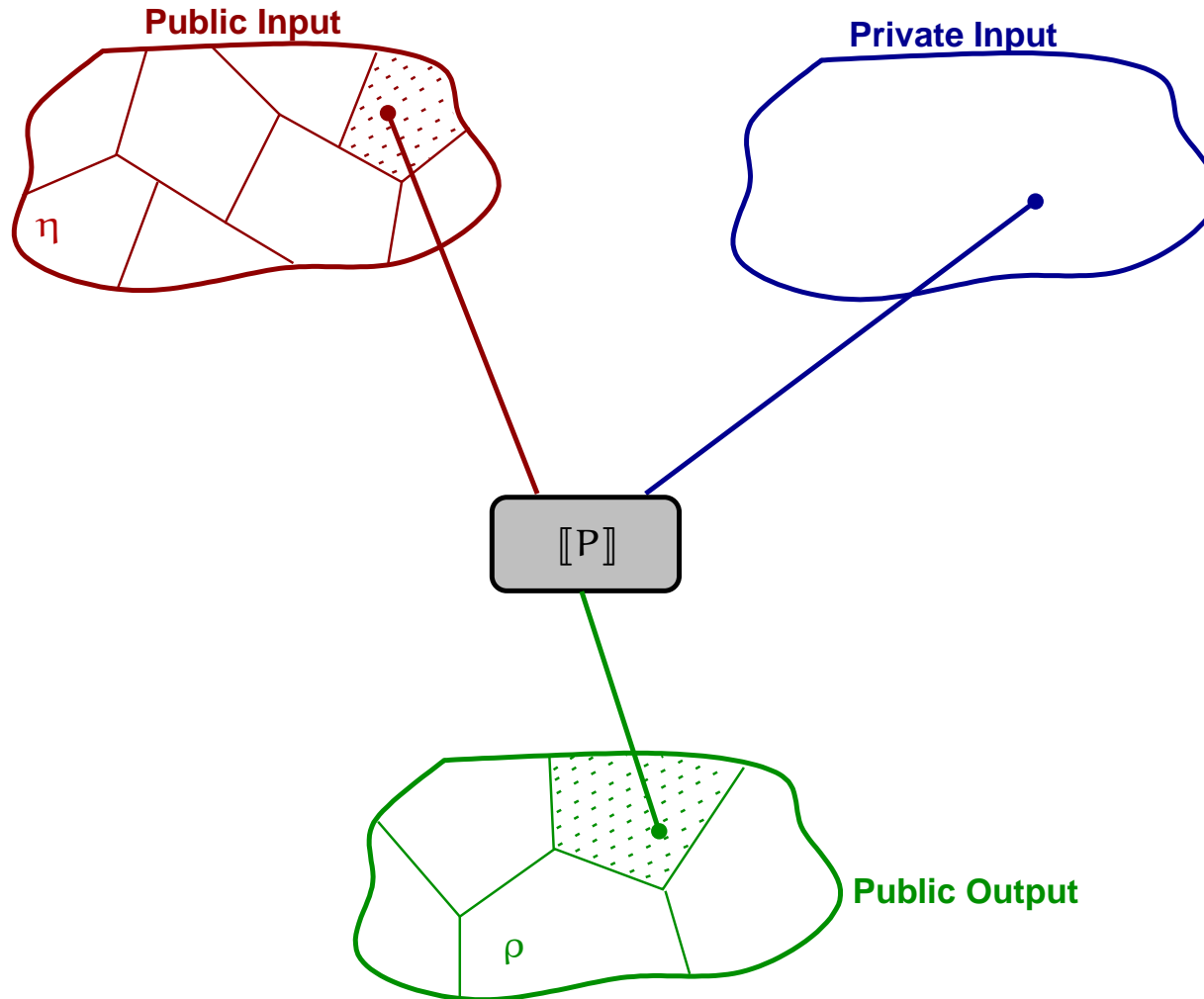
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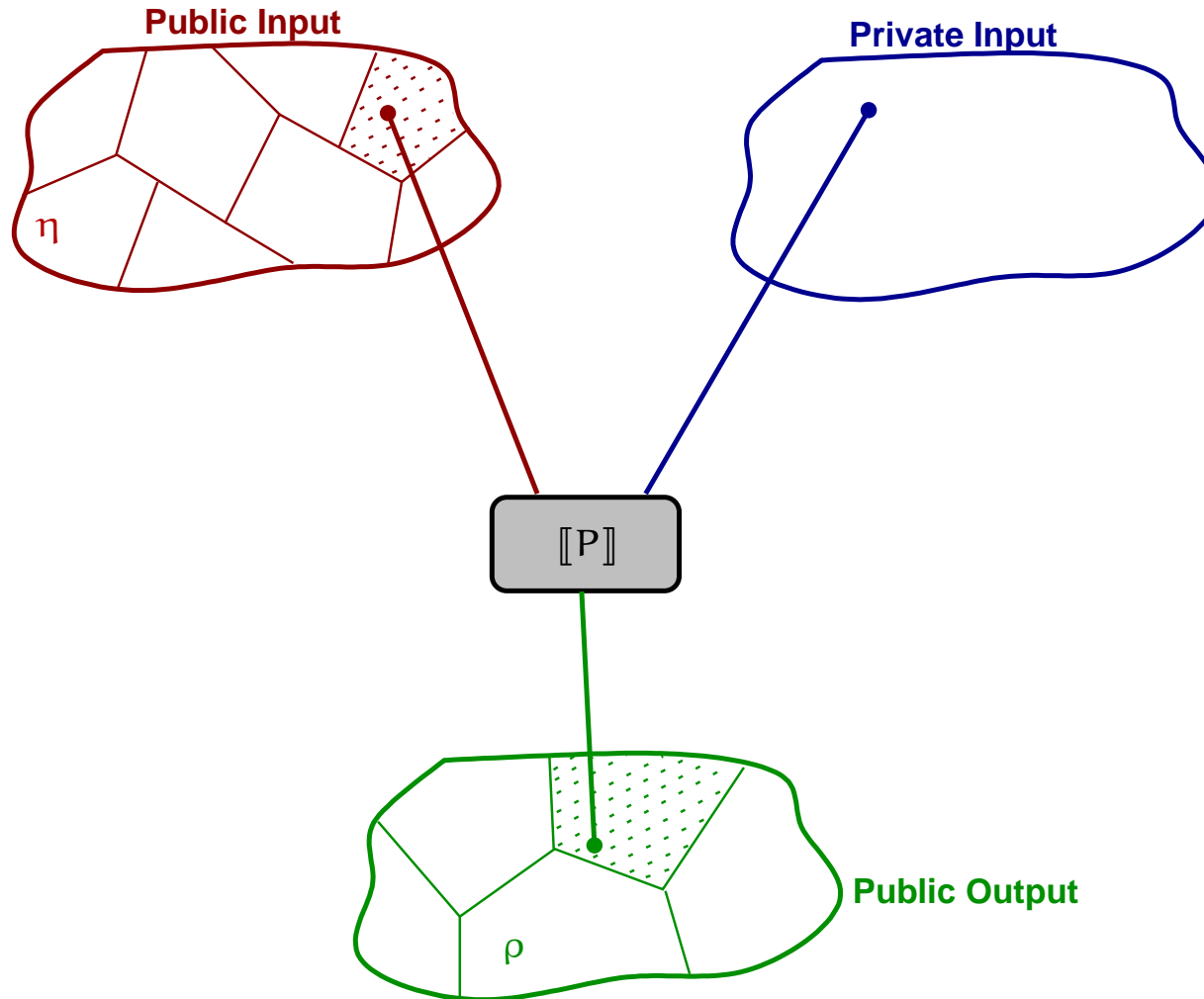
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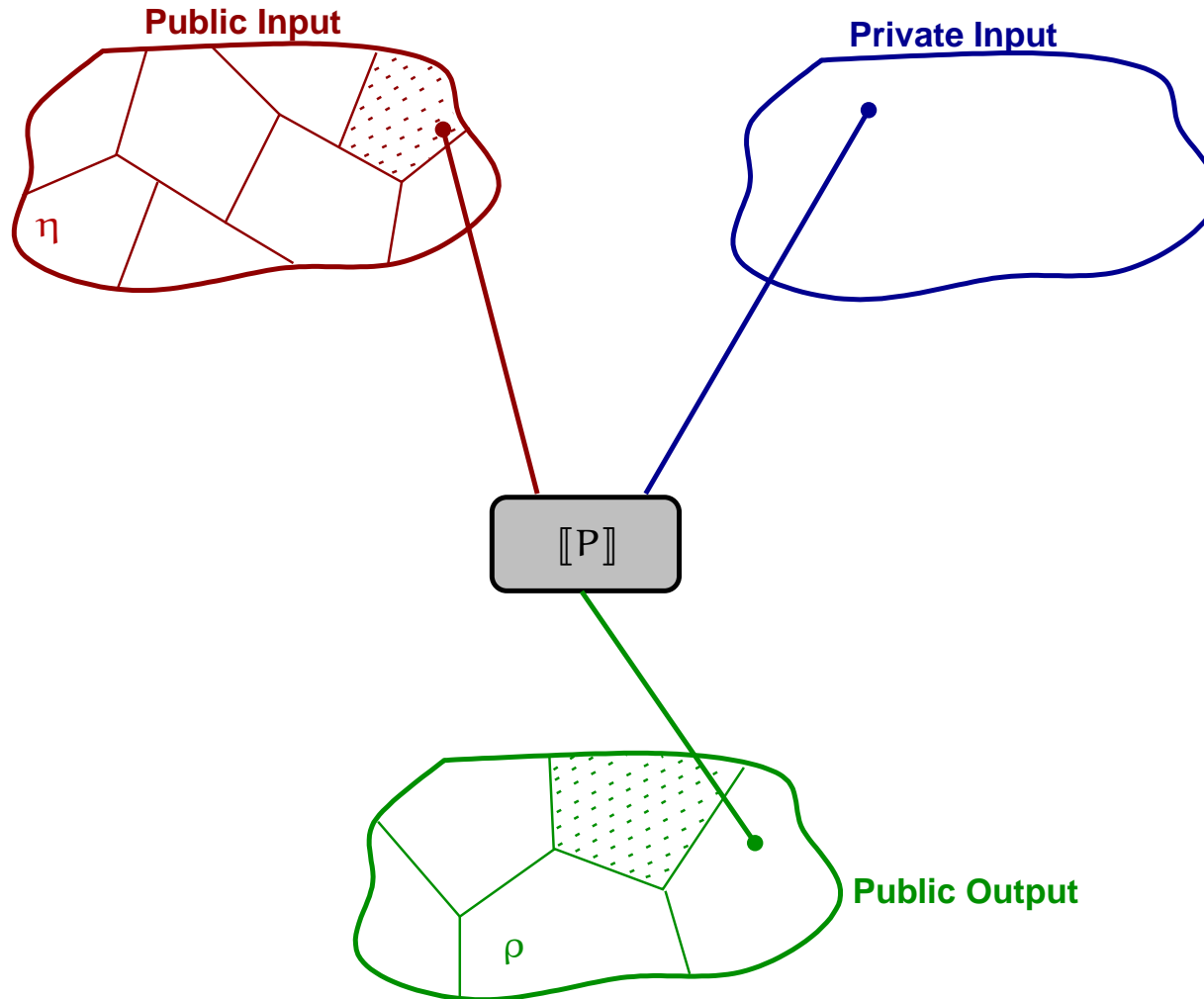
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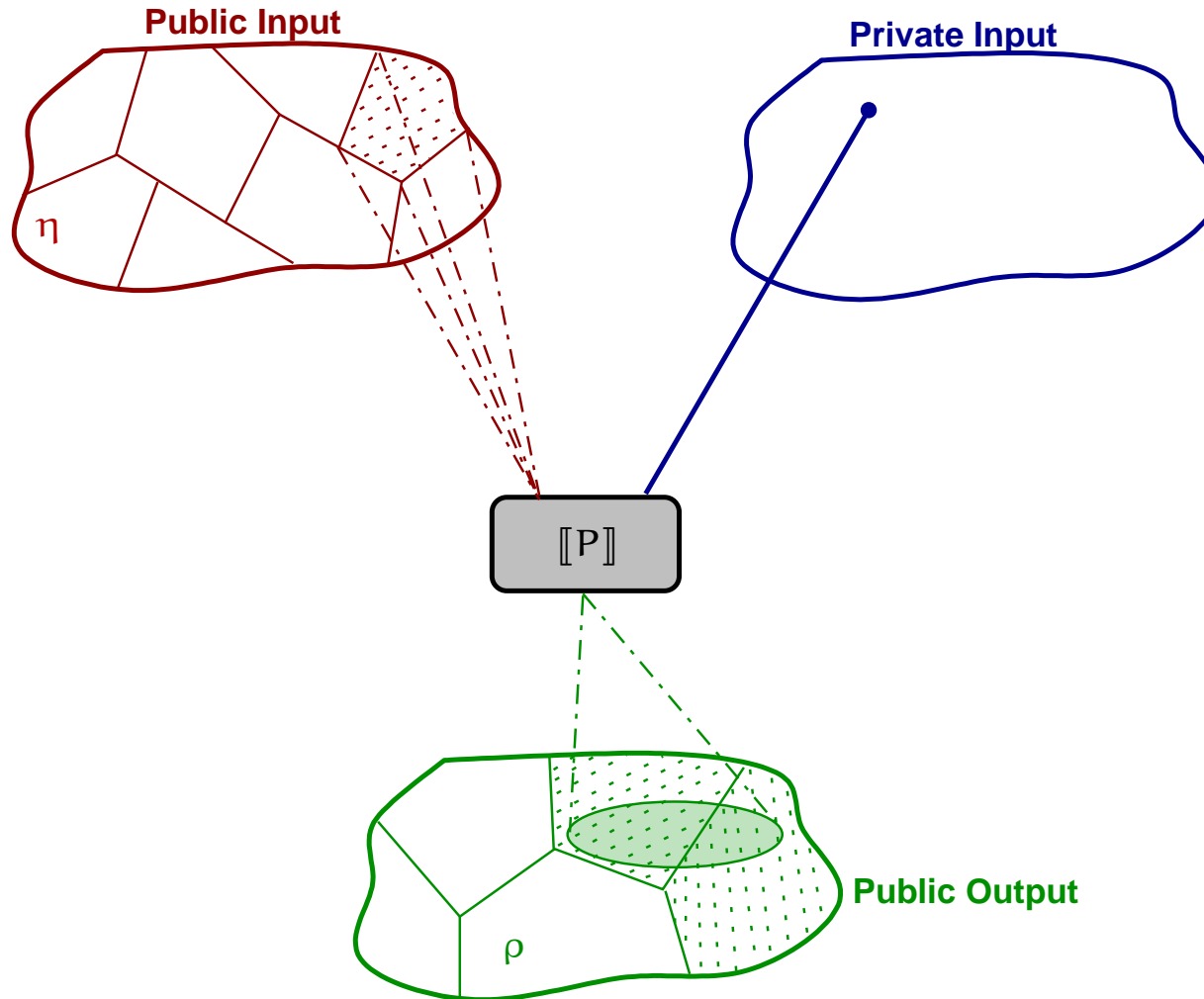
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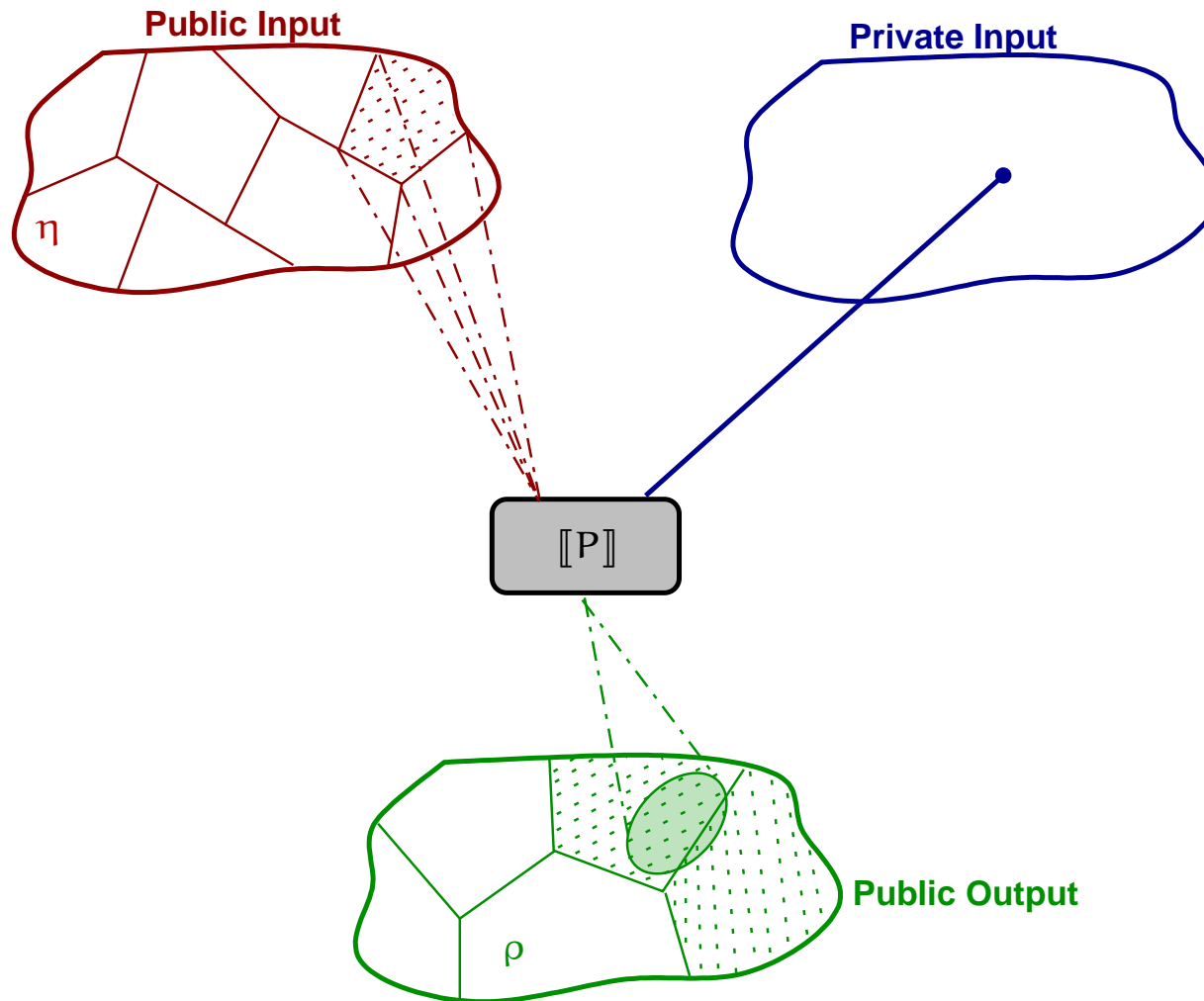
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$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)) : (\eta)P(\rho) :$$

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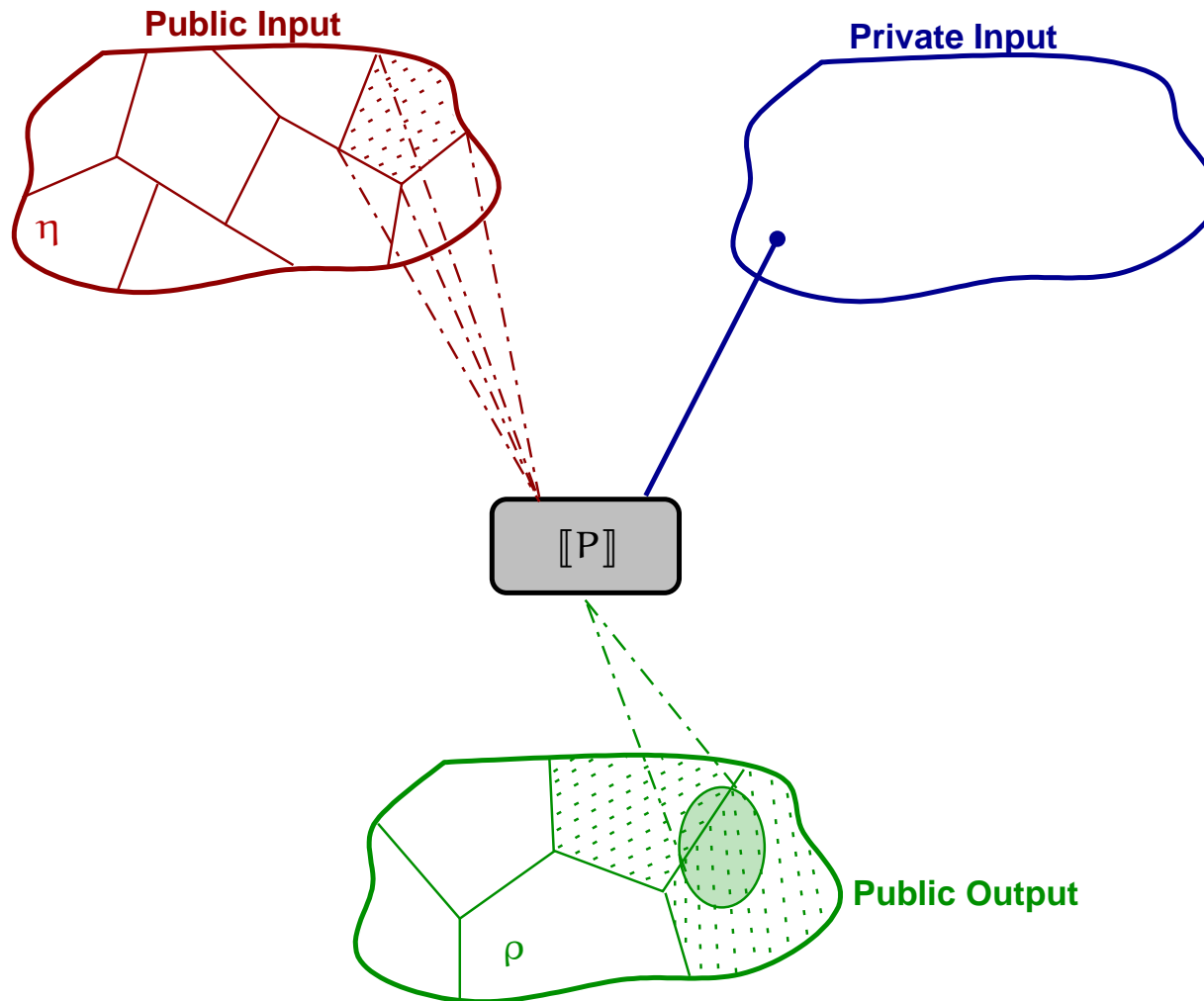
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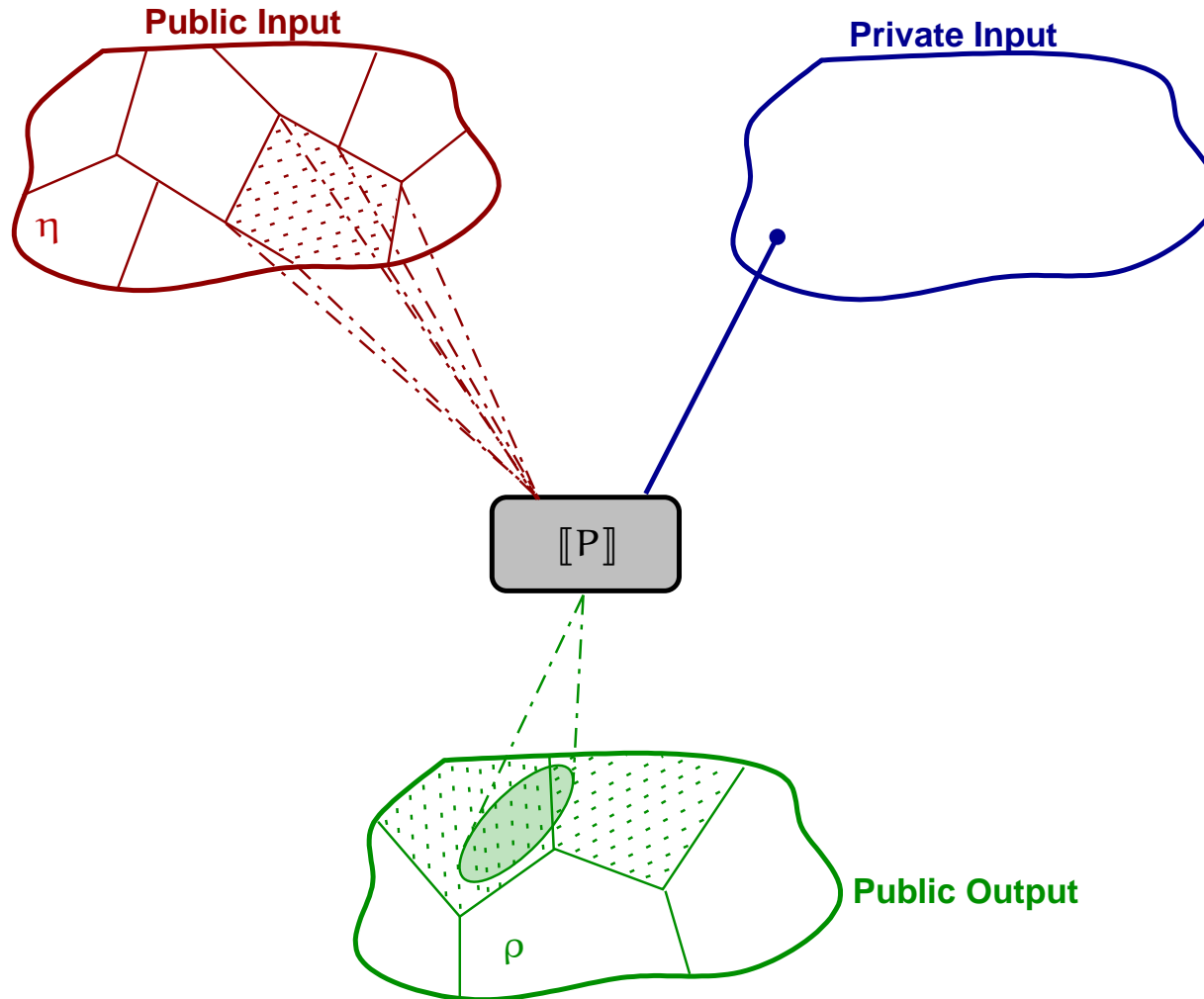
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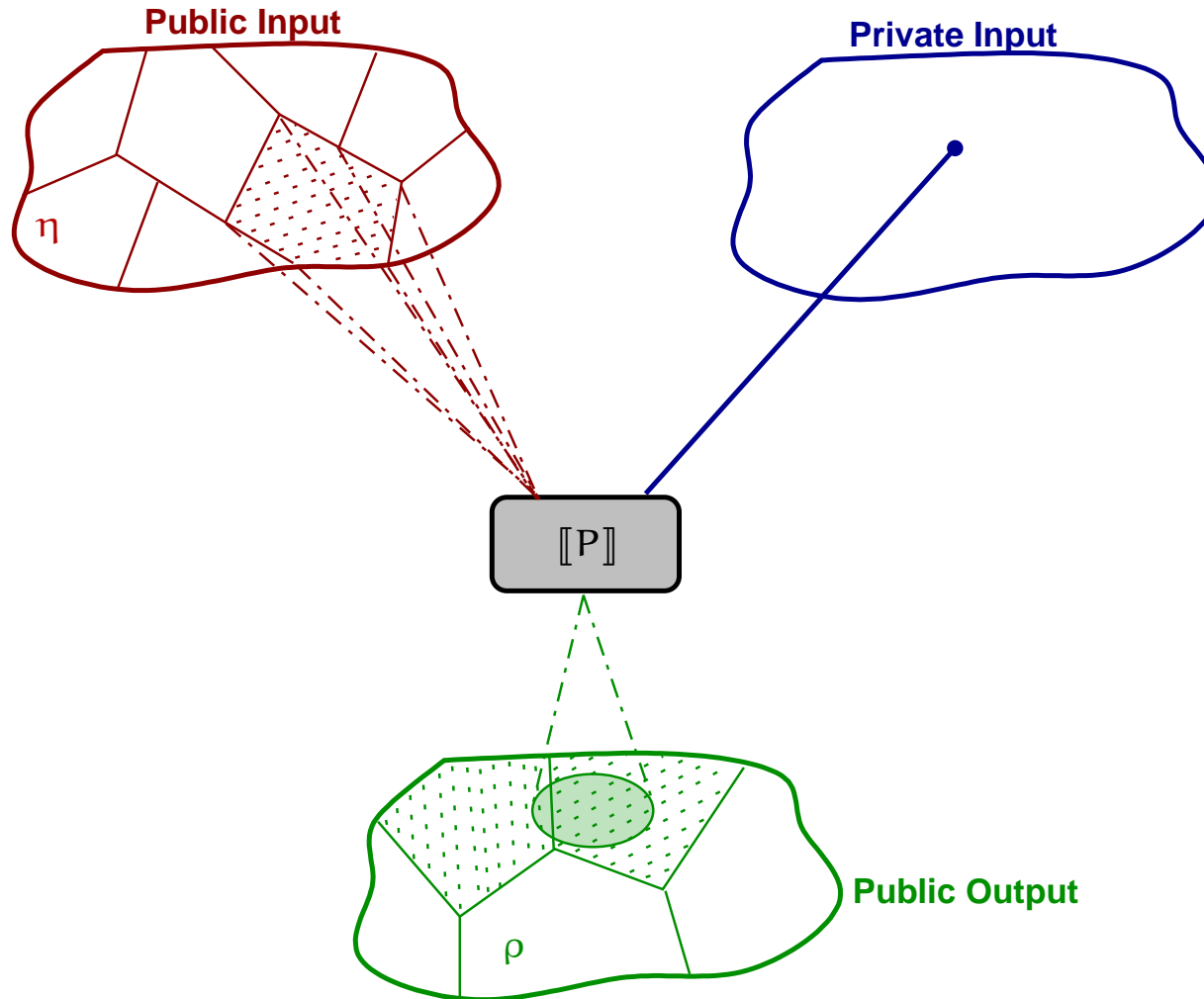
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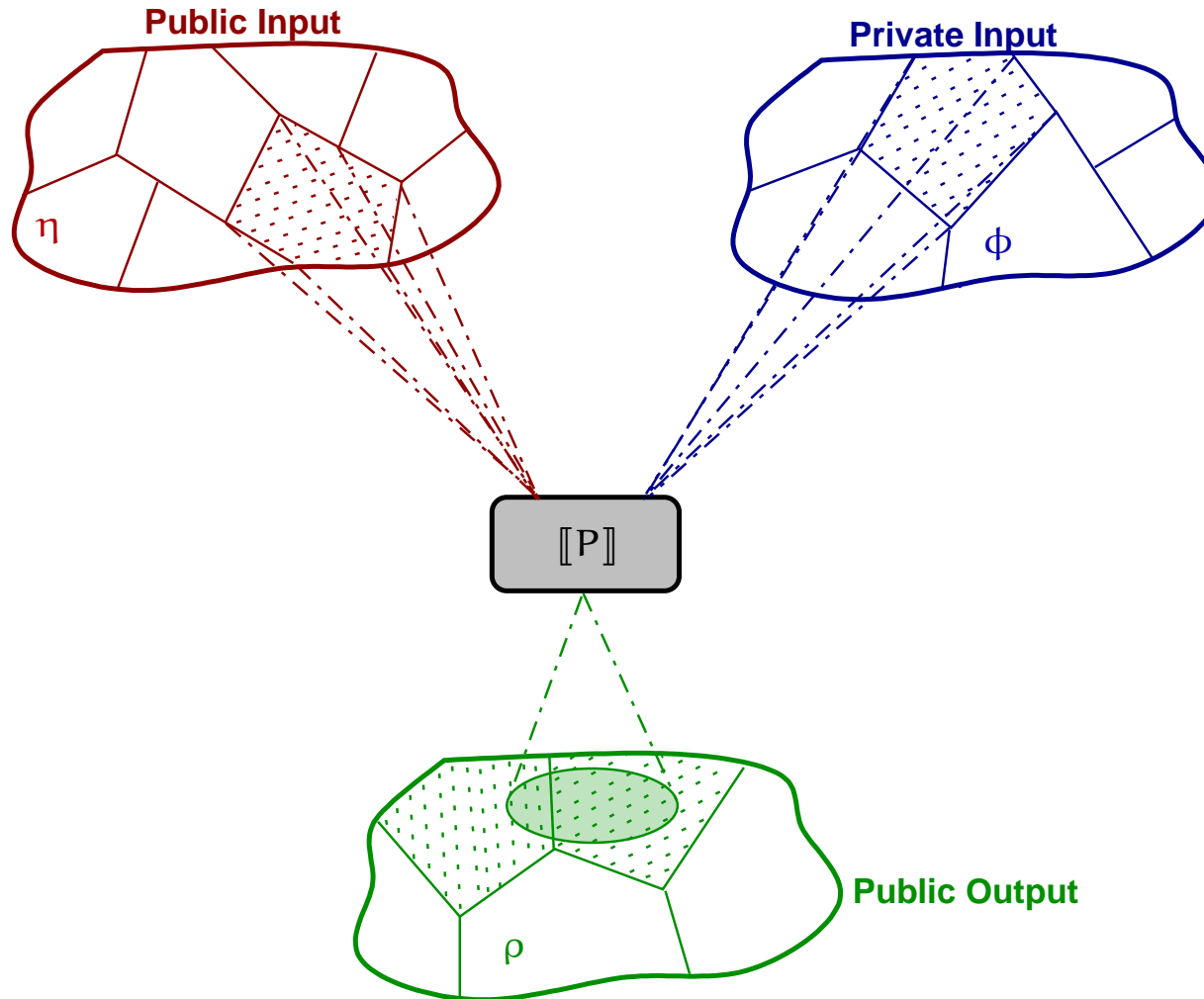
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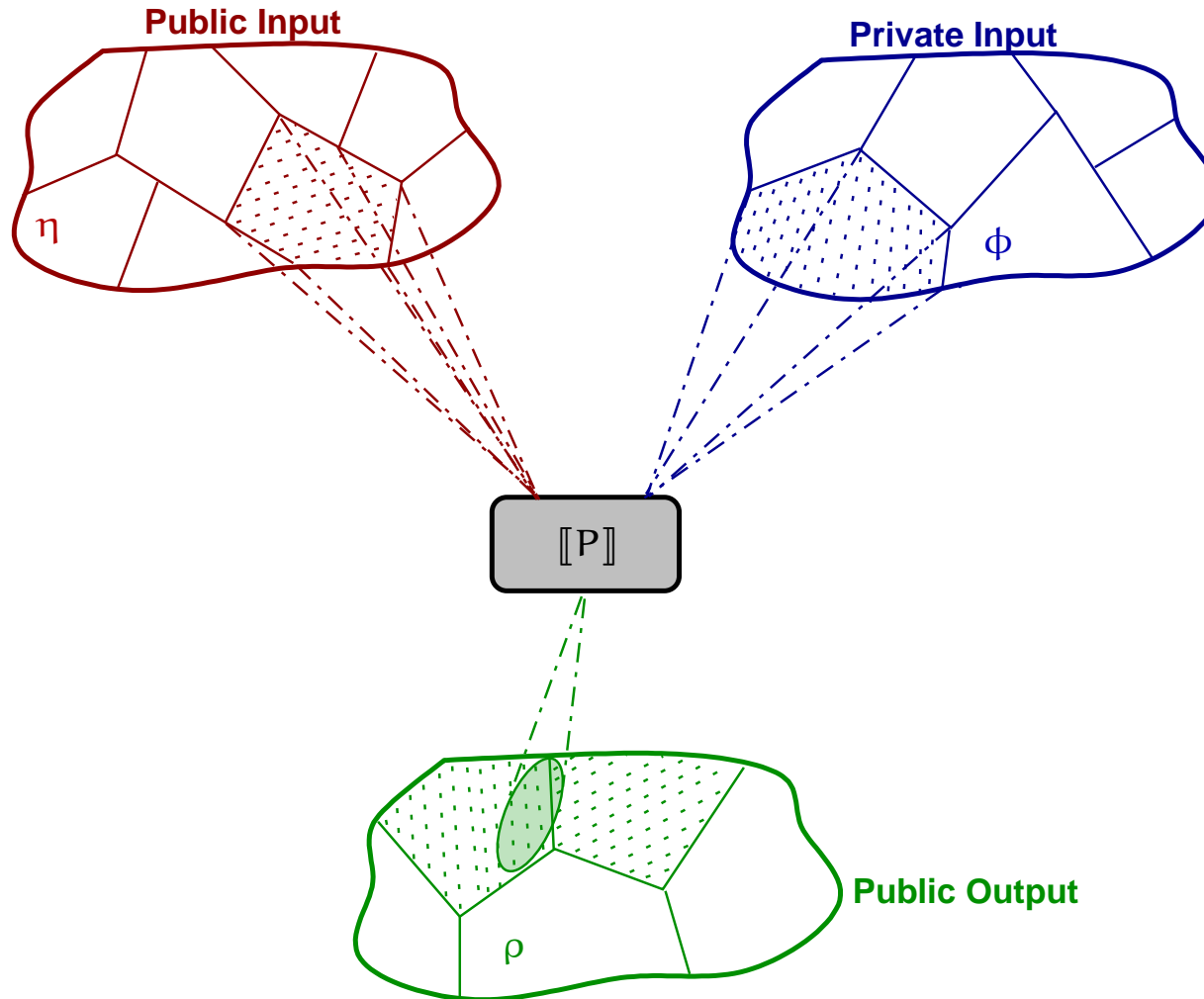
Abstracting non-interference II: ANI



$$\rho, \eta \in \text{Abs}(\wp(\mathbb{V}^L)), \phi \in \text{Abs}(\wp(\mathbb{V}^H)): (\eta)P(\phi \rightsquigarrow \rho):$$

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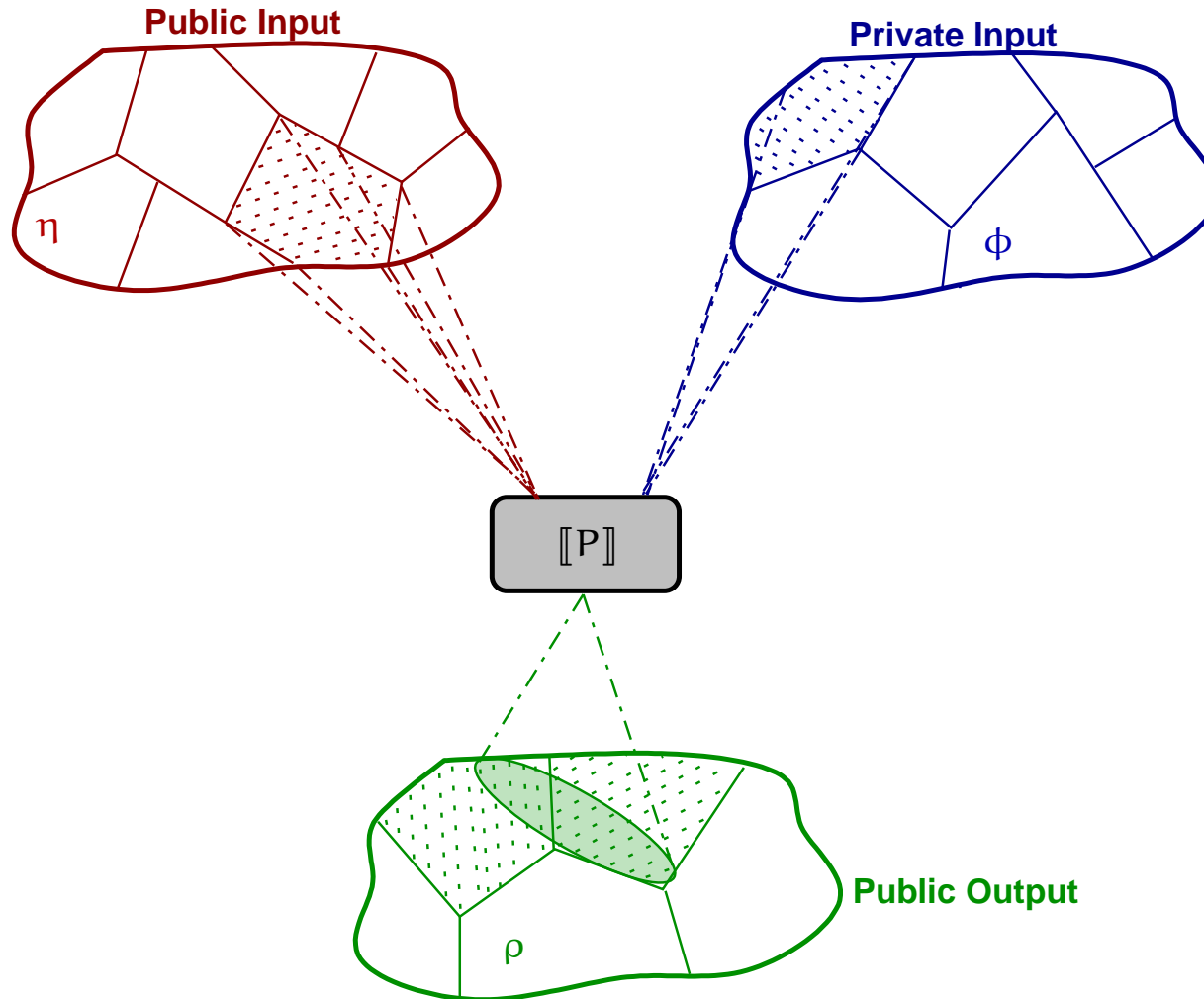
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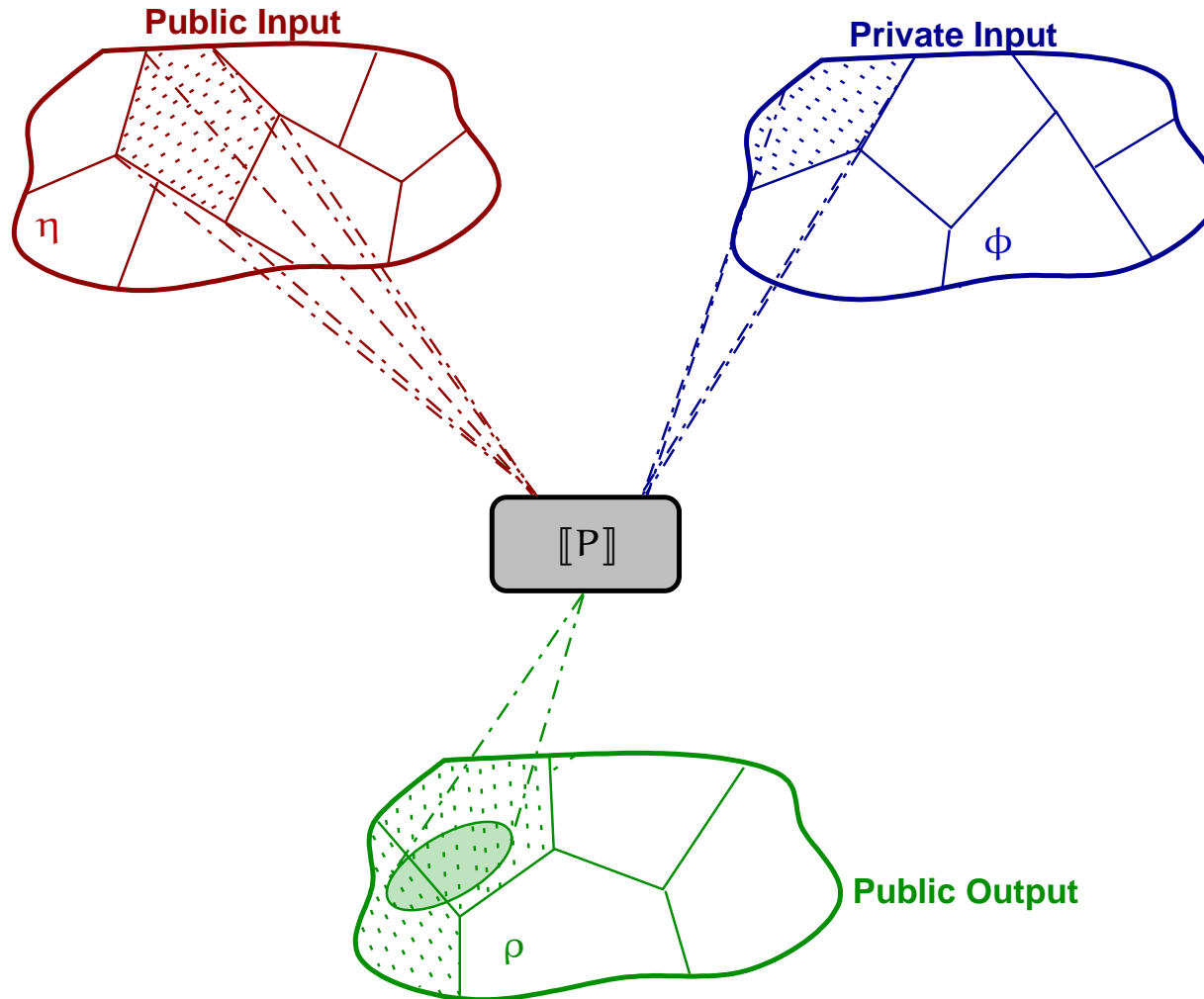
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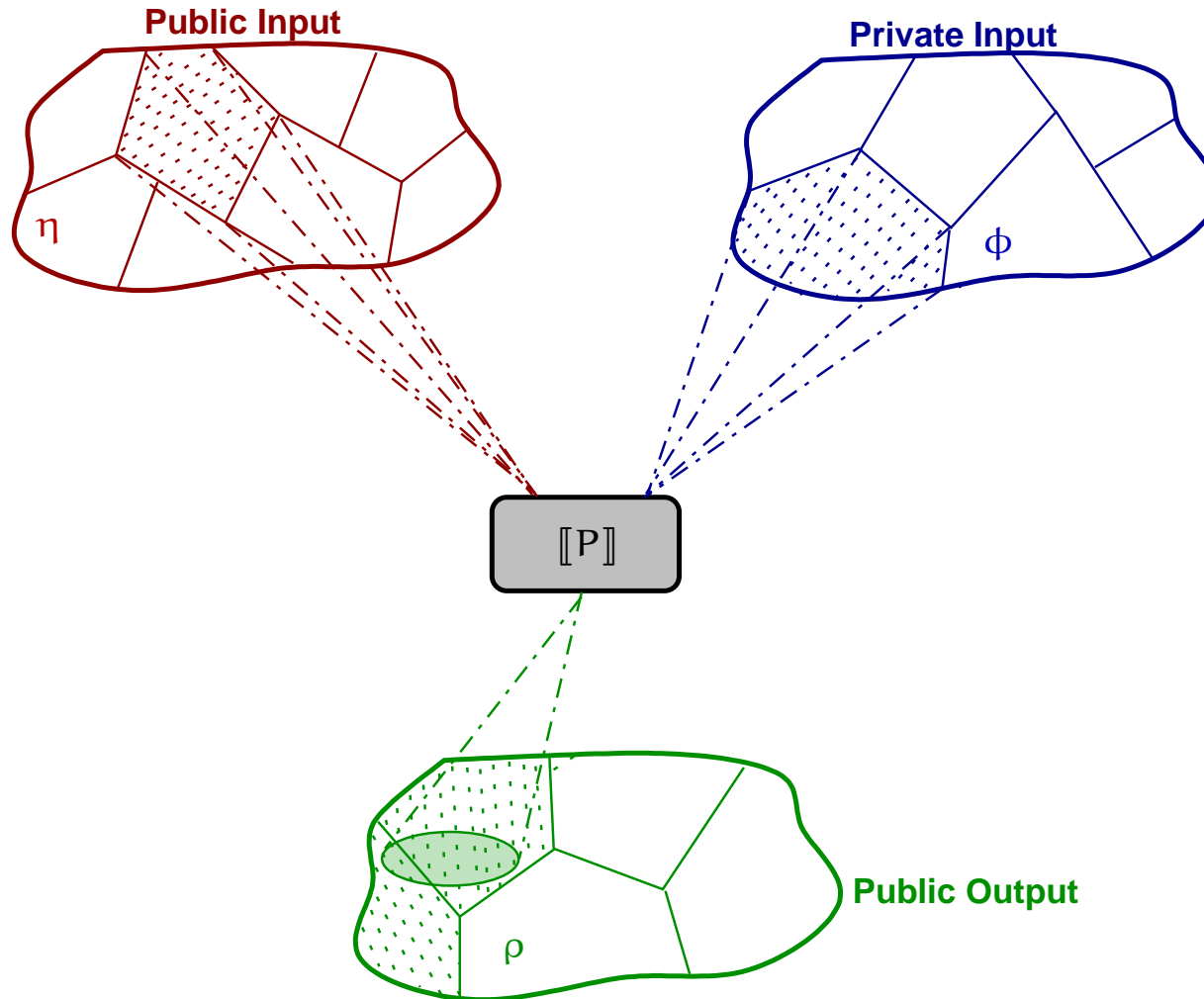
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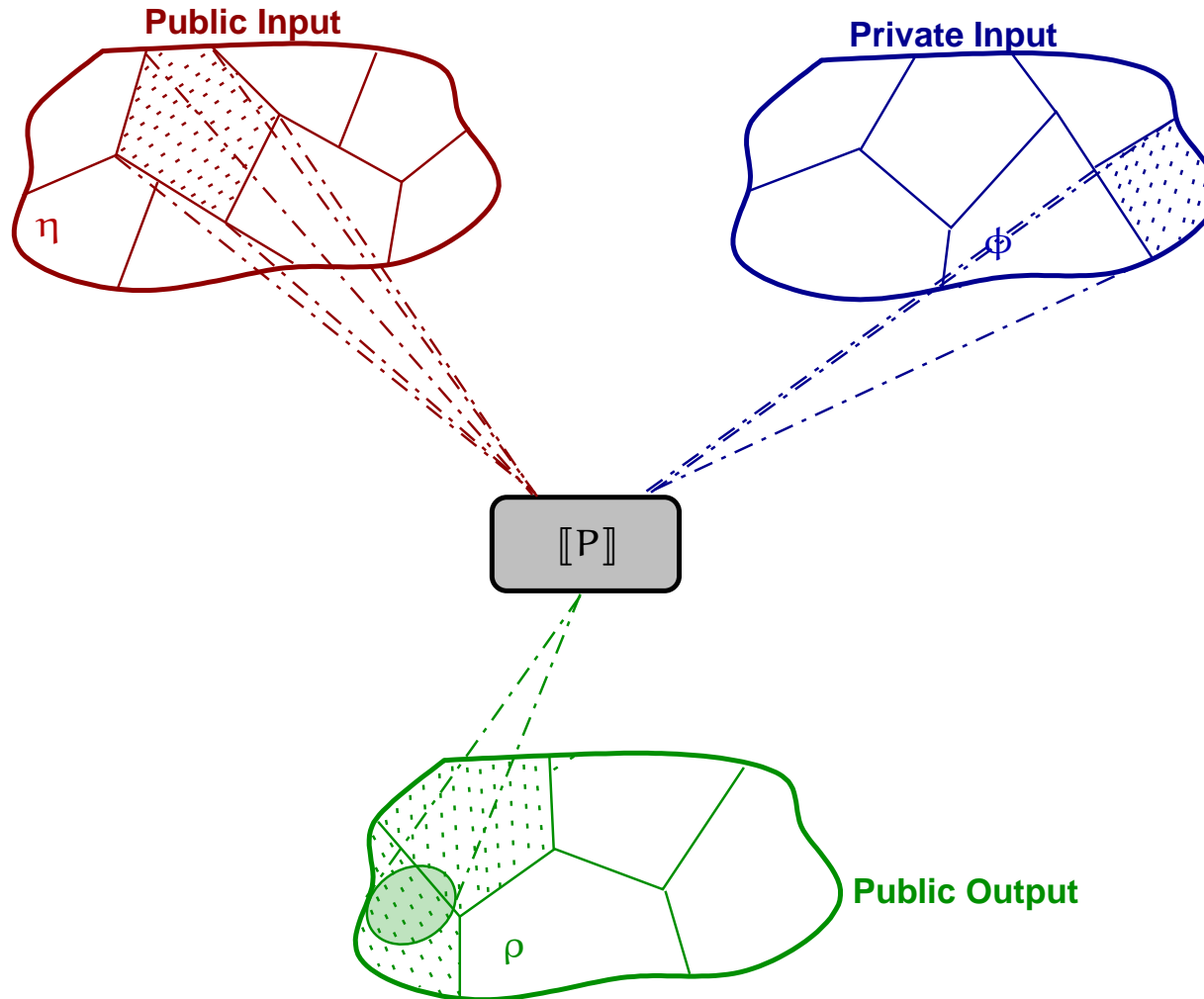
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Deriving output attackers

Abstract interpretation provides advanced methods for designing abstractions
(refinement, simplification, compression ...) [Giacobazzi & Ranzato '97]

Designing abstractions = designing attackers

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- ⑥ Characterize the most concrete ρ such that $(\eta)P(\phi \rightsquigarrow \rho)$
[The most powerful *public observer*]

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- ⑥ Characterize the most concrete ρ such that $(\eta)P(\phi \rightsquigarrow \rho)$
[The most powerful *public observer*]

⇒ This would provide a certificate for security with a fixed input observation.

Equivalence relation vs Closure Operators

Equivalence relations uniquely correspond to particular upper closure operators:

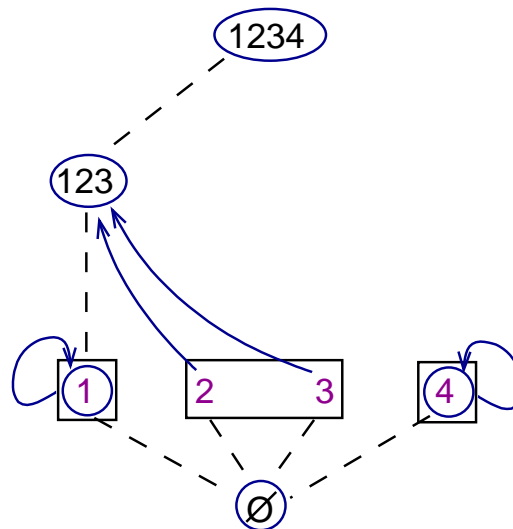
PARTITIONING
[Ranzato and Tapparo'04]

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$$\textcircled{G} \quad \eta \in \text{uco}(\wp(C)) \Rightarrow \forall x, y. x \mathit{Rel}^\eta y \text{ iff } \eta(x) = \eta(y)$$

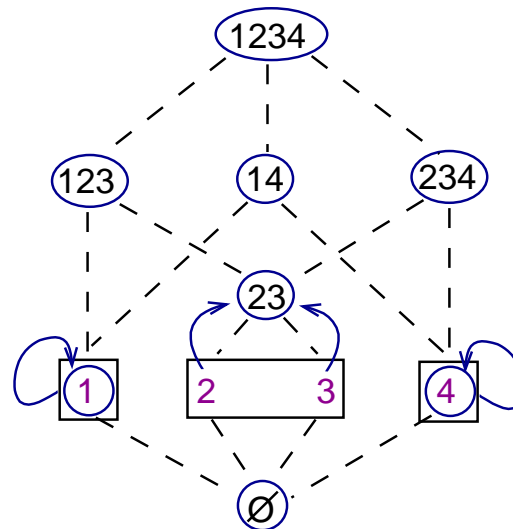


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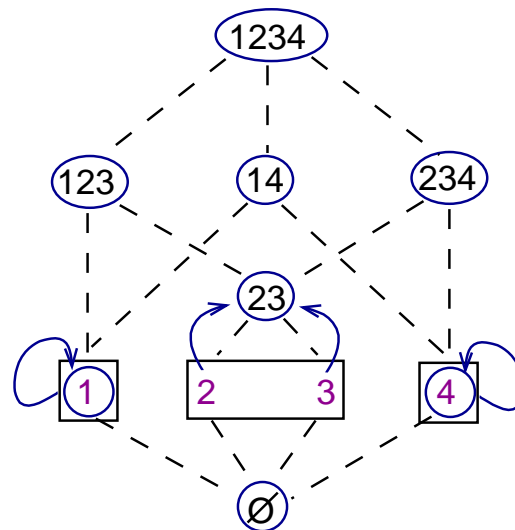


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- ⑥ $R \in Eq(C) \Rightarrow \mathit{Clo}^R(X) \stackrel{\text{def}}{=} \bigcup_{x \in X} [x]_R$



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$$\textcircled{6} \quad \Pi(\eta) \stackrel{\text{def}}{=} \mathbf{Clo}^{\mathbf{Rel}^\eta} \sqsubseteq \eta$$

PER model vs Narrow ANI

PROPOSITION:

$$\begin{aligned} [\eta]P(\rho) & \text{ iff } \llbracket P \rrbracket : All \times Rel^n \rightarrow All \times Rel^p \\ & \text{ iff } [\Pi(\eta)]P(\Pi(\rho)) \end{aligned}$$

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$$[R]P(S) \text{ iff } \llbracket P \rrbracket : All \times R \rightarrow All \times S$$

PER model vs ANI

PROPOSITION 2:

$$(\eta)P(\phi \rightsquigarrow \rho) \text{ if } (\Pi(\eta))P(\Pi(\phi) \rightsquigarrow \Pi(\rho))$$

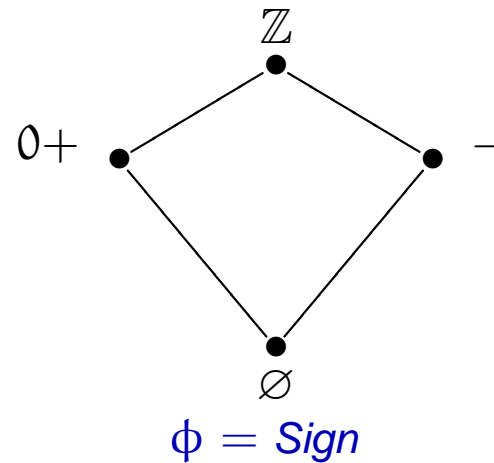
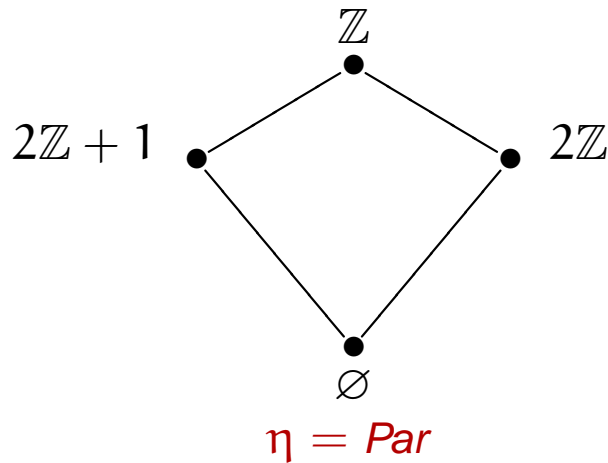
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EXAMPLE:

$P \stackrel{\text{def}}{=} \text{if } h = 0 \text{ then } l := l \bmod 6 + 2; \text{ else if } l < 0 \text{ then } l := 2 \text{ else } l := 7;$



$\rho = \text{Int}$

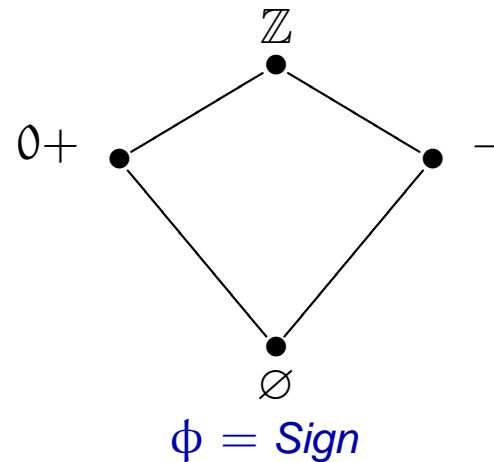
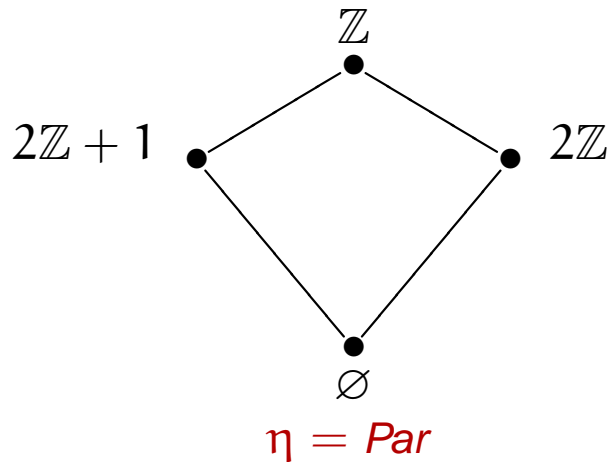
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$$\eta(l) = 2\mathbb{Z} : \quad \rho(\llbracket P \rrbracket(0+, 2\mathbb{Z})^\perp) = \rho(\{2, 3, 5, 7\}) = [2, 7] =$$

$$\rho(\llbracket P \rrbracket(-, 2\mathbb{Z})^\perp) = \rho(\{2, 7\})$$

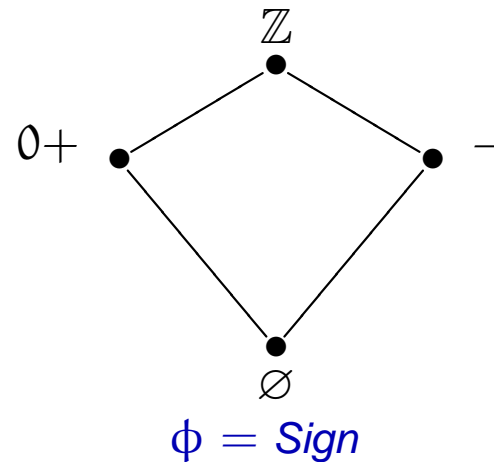
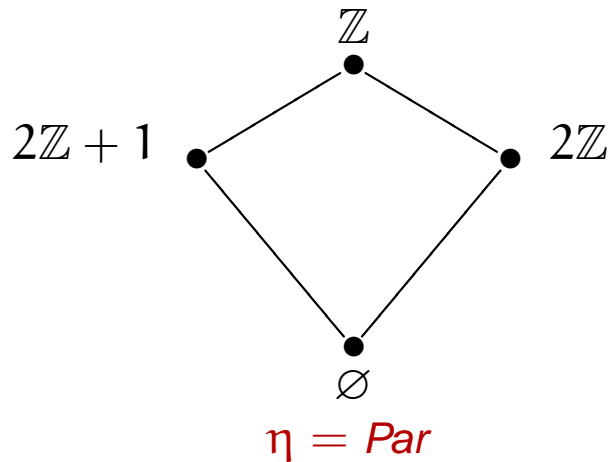
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Deriving un-constrained attackers

Input Observation: R ;

Output Observation: S ;

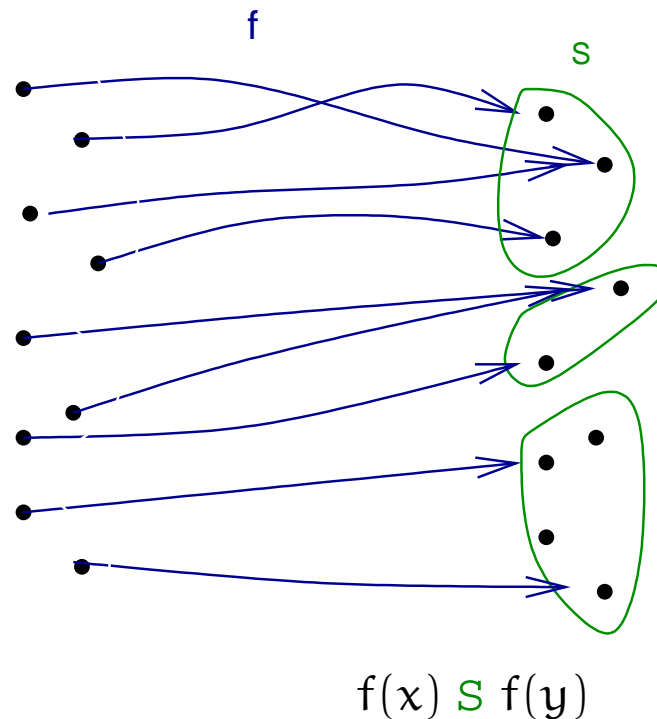
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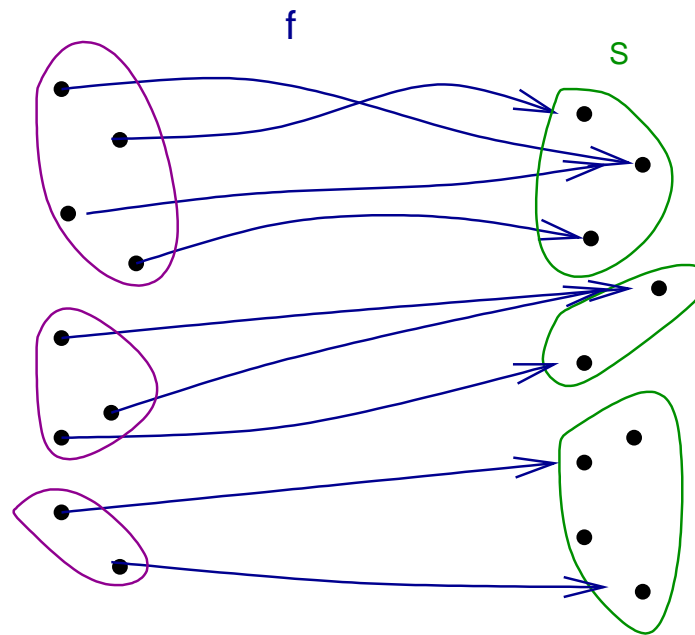


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$$x \hat{f}^{-1}(S) y \text{ iff } f(x) S f(y)$$

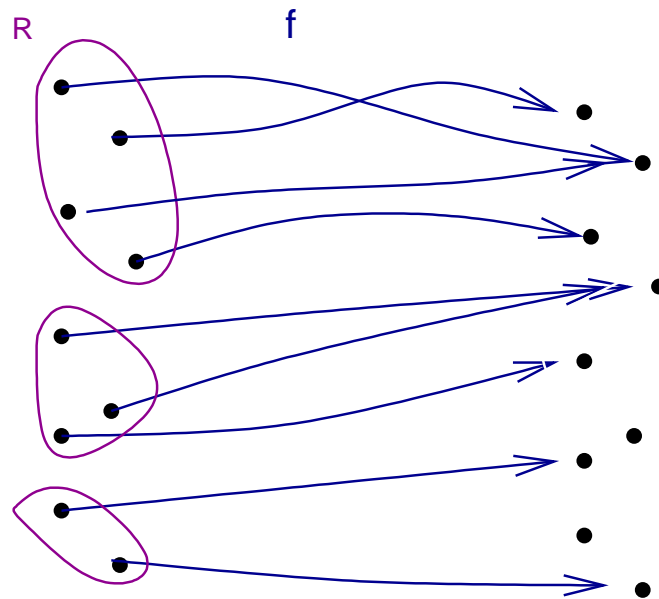
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$$x \hat{f}^{-1}(S) y \text{ iff } f(x) S f(y)$$



$$x' R y'$$

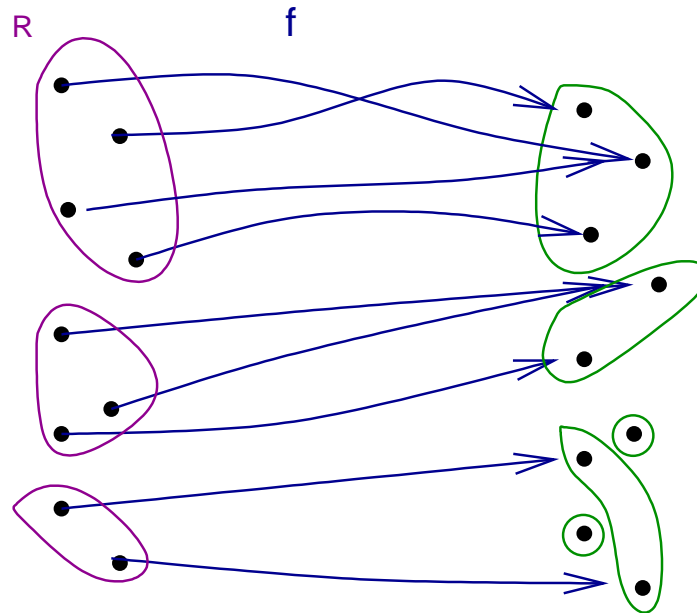
Deriving un-constrained attackers

Input Observation: R ;

Output Observation: S ;

Non-Interference: $f : R \rightarrow S$ iff $x R y \Rightarrow f(x) S f(y)$

$$x \hat{f}^{-1}(S) y \text{ iff } f(x) S f(y)$$



$$x \hat{f}(R) y \text{ iff } x = y \text{ or } x = f(x'), y = f(y'), x' R y'$$

Deriving un-constrained attackers

Input Observation: R ;

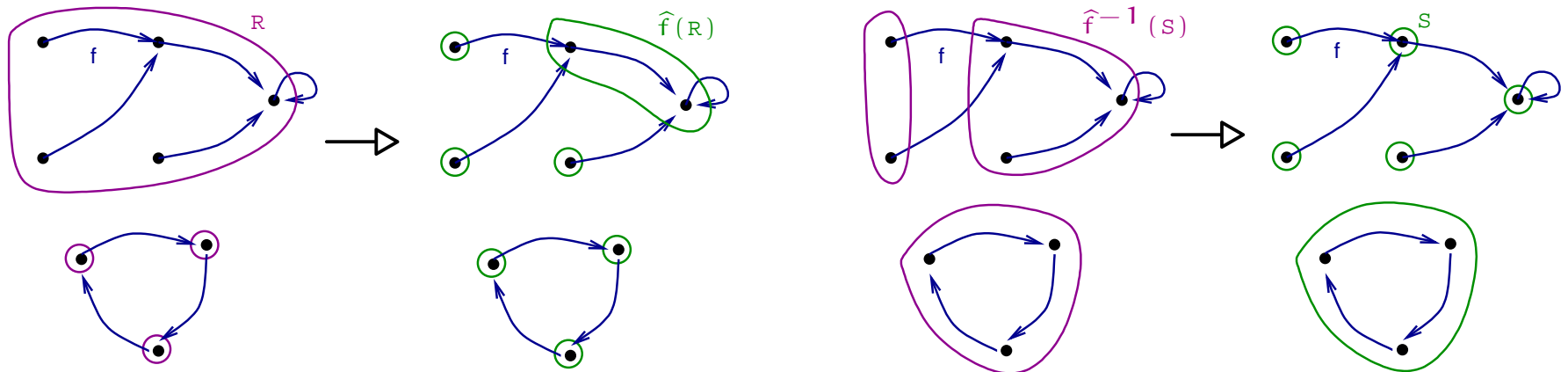
Output Observation: S ;

Non-Interference: $f : R \rightarrow S$ iff $x R y \Rightarrow f(x) S f(y)$

$$x \hat{f}^{-1}(S) y \text{ iff } f(x) S f(y)$$

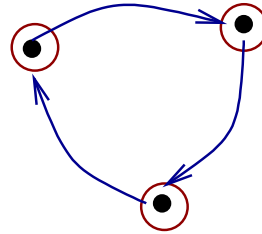
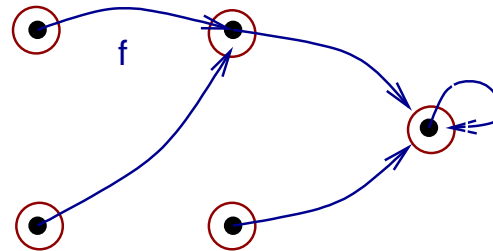
$$x \hat{f}(R) y \text{ iff } x = y \text{ or } x = f(x'), y = f(y'), x' R y'$$

EXAMPLE:



Fix point attackers

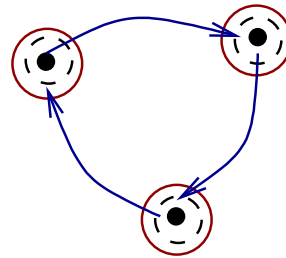
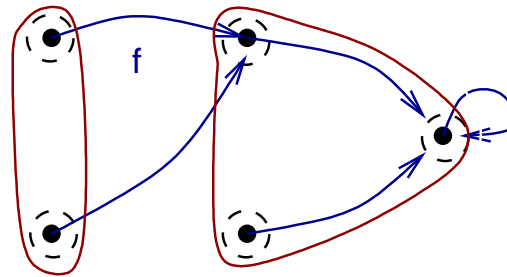
⑥ Least fix point of \hat{f}^{-1}



(a)

Fix point attackers

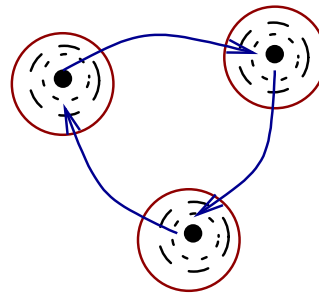
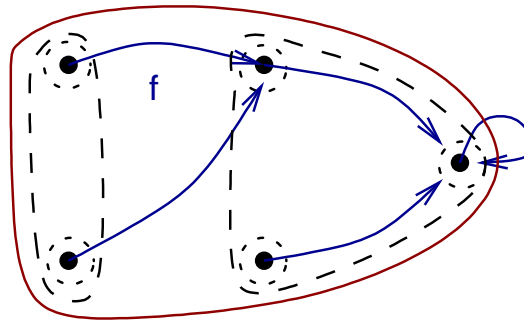
⑥ Least fix point of \hat{f}^{-1}



(a)

Fix point attackers

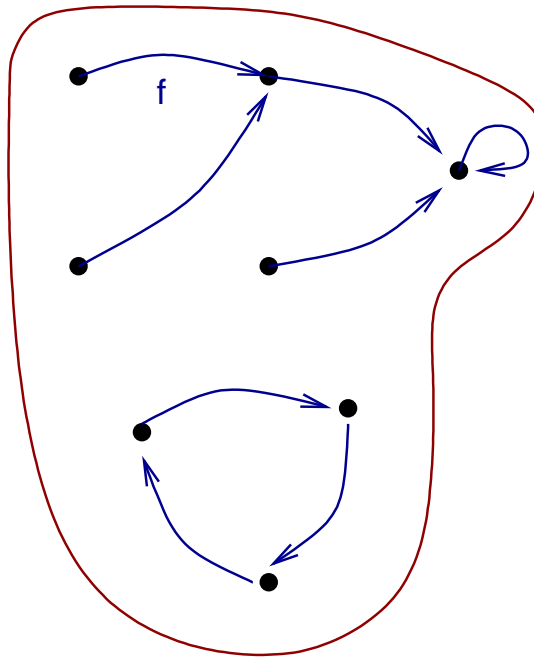
⑥ Least fix point of \hat{f}^{-1}



(a)

Fix point attackers

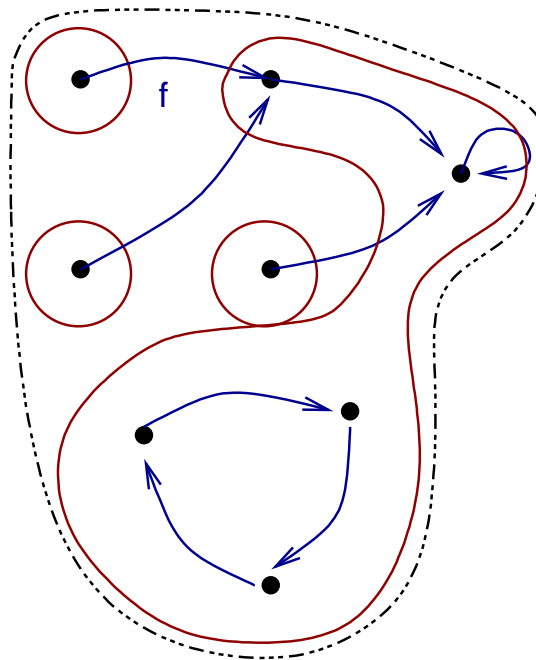
- ⑥ Least fix point of \hat{f}^{-1}
- ⑥ Greatest fix point of \hat{f}



(b)

Fix point attackers

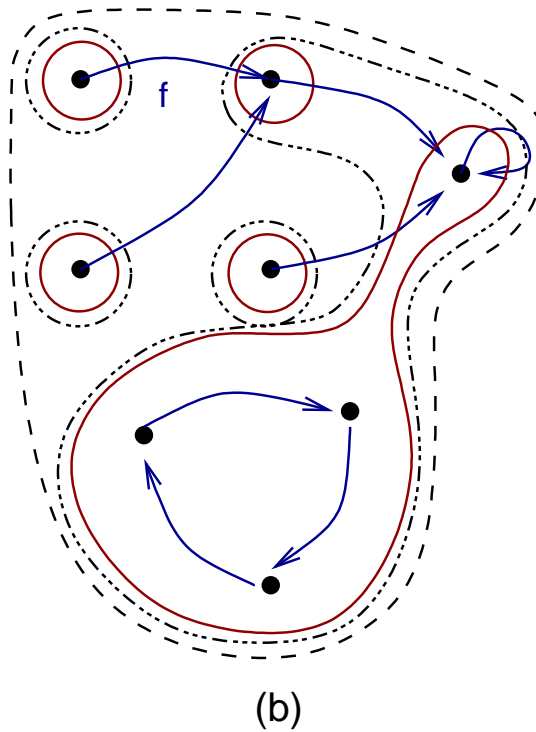
- ⑥ Least fix point of \hat{f}^{-1}
- ⑥ Greatest fix point of \hat{f}



(b)

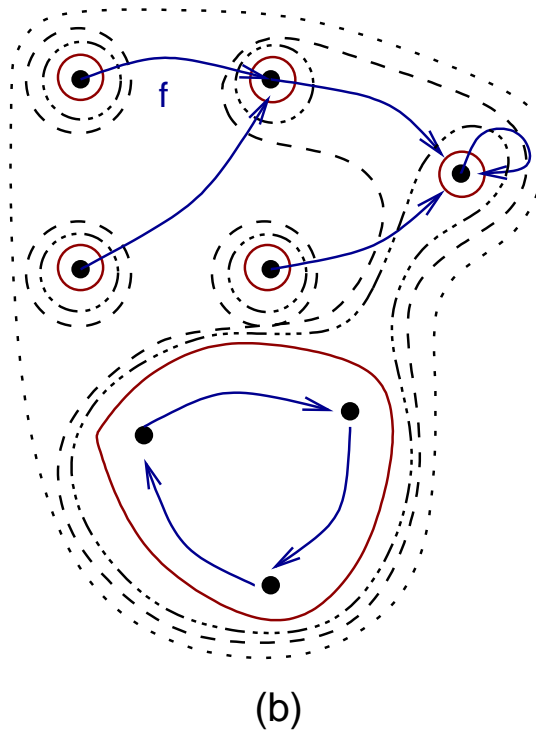
Fix point attackers

- ⑥ Least fix point of \hat{f}^{-1}
- ⑥ Greatest fix point of \hat{f}



Fix point attackers

- ⑥ Least fix point of \hat{f}^{-1}
- ⑥ Greatest fix point of \hat{f}



Fix point attackers

- ⑥ Least fix point of \hat{f}^{-1}
- ⑥ Greatest fix point of \hat{f}
- ⑥ The smallest \mathbb{R} such that $[\mathbb{R}]P(\mathbb{R})$ is the least fix point of $\hat{f}(A// \times \mathbb{R})$.

PER model and completeness

Completeness in abstract interpretation can be formalized as follows:

$$\rho \circ f \circ \eta = \rho \circ f$$

[Giacobazzi et al. '00]

iff

$$\forall x, y. \eta(x) = \eta(y) \Rightarrow \rho(f(x)) = \rho(f(y))$$

PER model and completeness

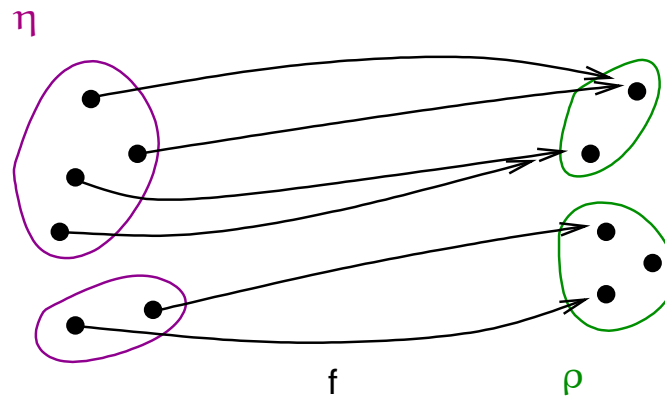
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PER model and completeness

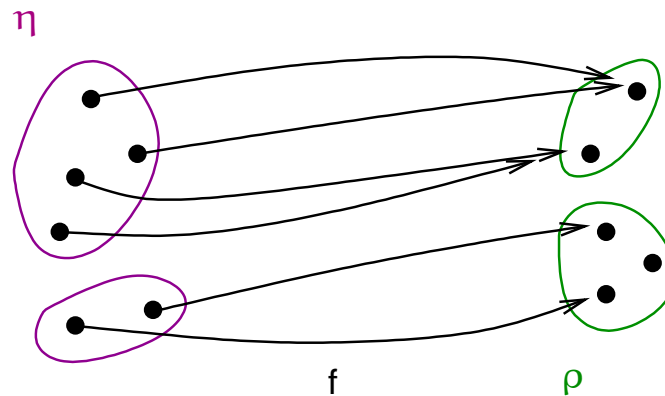
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$$f : R \rightarrow S$$

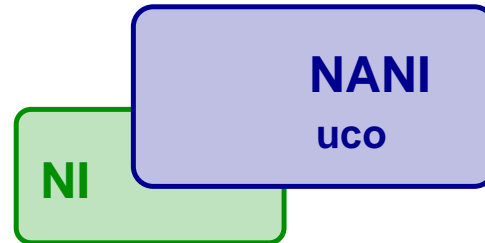
iff

$$Clo^S \circ f \circ Clo^R = Clo^S \circ f$$

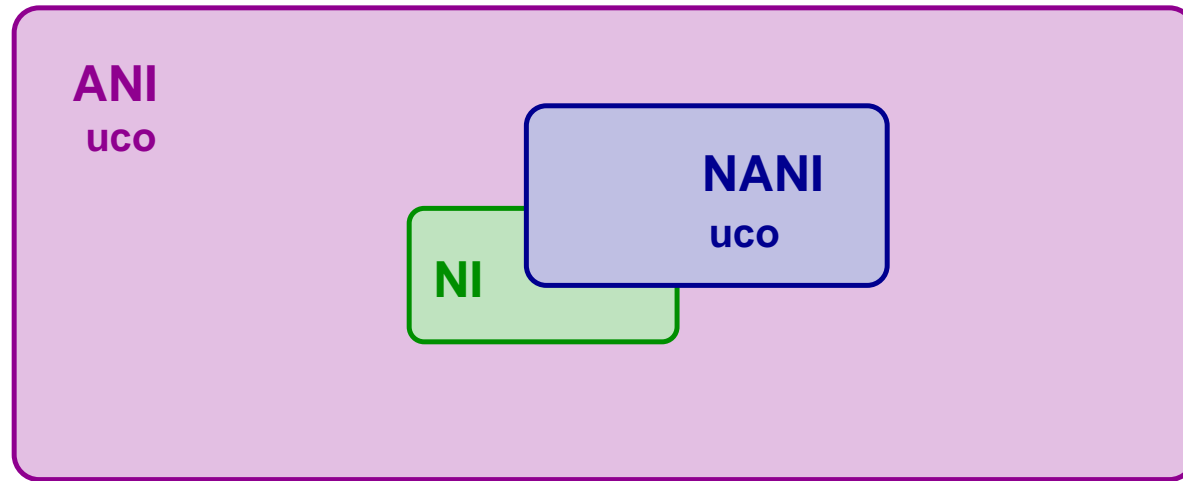
Conclusion

NI

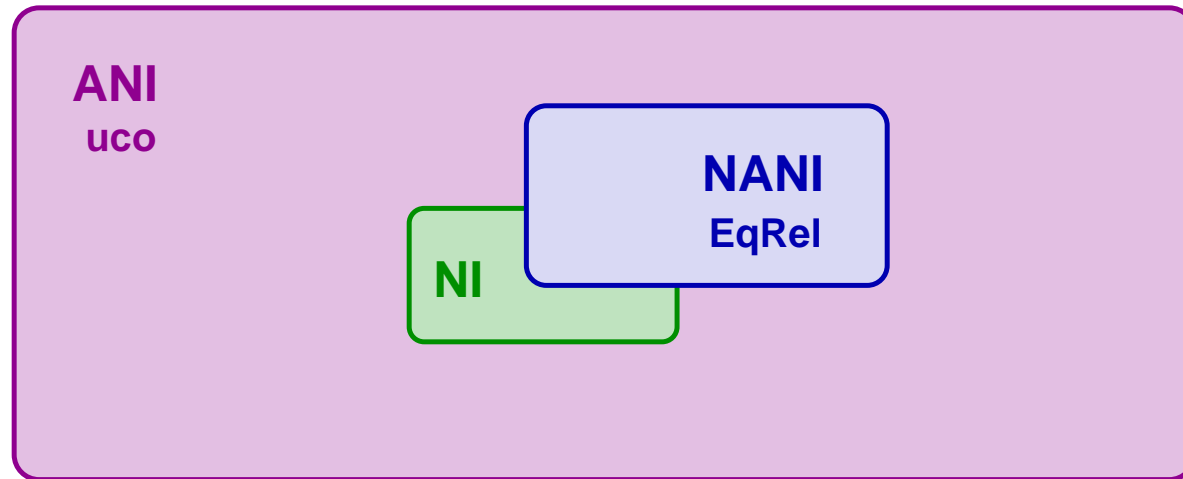
Conclusion



Conclusion



Conclusion



Conclusion

