

Normal, Abby Normal, Prefix Normal

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Prefix normal words

Fici, Lipták (DLT 2011)

Definition

A binary word w is **prefix normal** (w.r.t. 1) if no substring has more 1s than the prefix of the same length.

Example

$$w = 10110001001101110010$$

$$w' = 11101001011001010010$$

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Example

$w = 10110001001101110010$ *NO*

$w' = 11101001011001010010$ *YES*

all p.n. words of length 4:

1111, 1110, 1101, 1100, 1010, 1001, 1000, 0000.

Prefix normal games

Let's play a game.

Alice and Bob are constructing a binary string of length n together.

Alice wins if the word is prefix normal.

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|----|--------------|---|---|---|---|---|
| 1. | <i>start</i> | - | - | - | - | - |
| 2. | <i>Alice</i> | 1 | - | - | - | - |

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| 2. | <i>Alice</i> | 1 | - | - | - | - |
| 3. | <i>Bob</i> | 1 | 0 | - | - | - |

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Alice has won: 10000, 10100, 10101 are all p.n.

Prefix normal games

Exercise

Find n_0 s.t.

- Alice has a winning strategy for all $n < n_0$, and
- Bob has a winning strategy for all $n \geq n_0$.

Some properties of prefix normal words

- w is p.n. $\Rightarrow w0$ is p.n.
- w is p.n. \Rightarrow every prefix of w is p.n.
- for any w of length n , $v = 1^n w$ is p.n.
- ...

Prefix normal forms

- Let $F(w, k) = \max \#1\text{s in a substring of length } k$

Ex. $w = \mathbf{11010010101}$, $F(w, 1) = 1, F(w, 2) = 2, F(w, 3) = 2, \dots$

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- prefix normal form:**

For every w exists unique p.n. word w' s.t.

$$\forall k : F(w, k) = F(w', k)$$

$w' = \text{PNF}_1(w)$. (next slides)

Where do prefix normal words come from?

Binary Jumbled Pattern Matching

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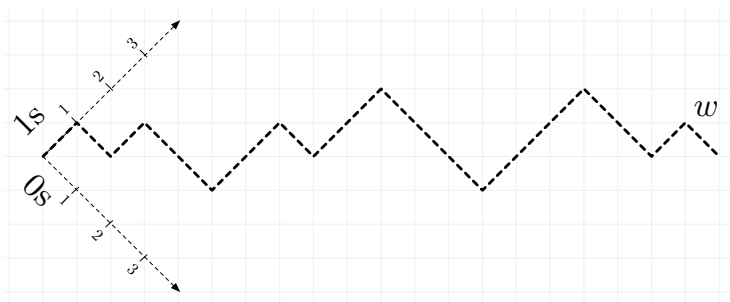
Recent papers on this problem:

PSC 2009, FUN 2010, IPL 2010, JDA 2012, ToCS 2012, IJFCS 2012, CPM 2012, SPIRE 2012, IPL 2013, IPL 2013, ESA 2013 \times 2, SPIRE 2013, TCS 2014, PhTRS-A 2014, CPM 2014, CPM 2014, ISIT 2014, ICALP 2014, ...

(red ones have an intersection with the authors of this paper)

BJPM with prefix normal forms

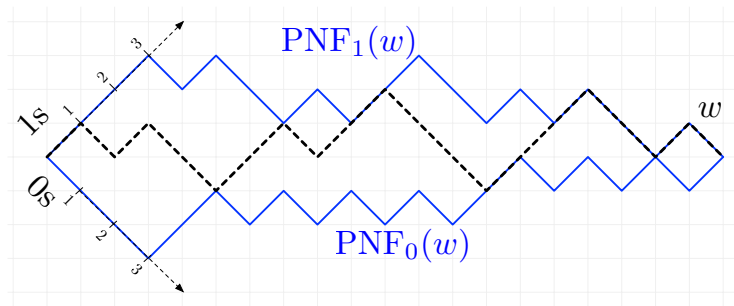
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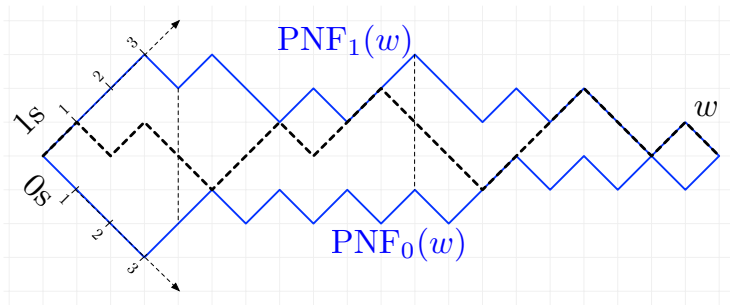
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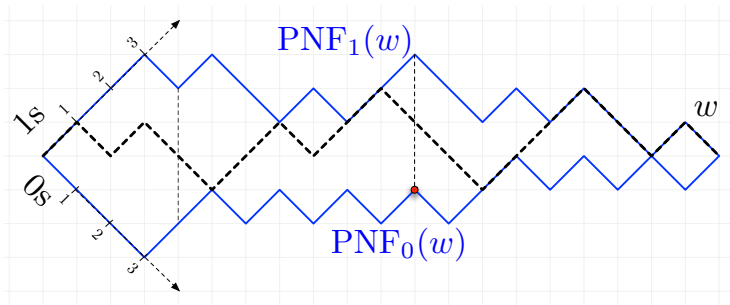
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BJPM with prefix normal forms

Does $w = \mathbf{10100110110001110010}$ have a substring of length 11 containing exactly 5 ones? **YES**



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BJPM with prefix normal forms

Does $w = \mathbf{10100110110001110010}$ have a substring of length 11 containing exactly 5 ones?

1s in $\text{pref}(\text{PNF}_1(w), 11) = 7$

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$$7 \geq 5 \geq 5 \rightsquigarrow \text{YES}$$

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Or might just be fun.

Some questions about prefix normal words

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 $O(n^2/\text{polylog } n)$, this paper: $O(n^2)$ w-c and $O(n)$ expected
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5. In general, **properties** of p.n. words.
– ongoing

Computing the PNF

Our contribution to **mechanical algorithm design**. Uses a folding ruler and Lipari's Canneto black sandy beach:



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Analysis

Uses a quadratic amount of sand but can be faster than other algorithms if implemented by a **very fast** person.

Enumerating p.n. words

Let $pnw(n) = \#$ prefix normal words of length n .

Easy:

$pnw(n)$ grows exponentially.

Proof: $\forall w : 1^{|w|}w$ is p.n. $\Rightarrow pnw(2n) \geq 2^n$.

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Theorem

Exists $c > 0$: $pnw(n) = \Omega(2^{c\sqrt{n \ln n}}) = \Omega((2 - \epsilon)^n) \quad \forall \epsilon > 0$.

(Proof uses a variant of the prefix normal game.)

Theorem

$pnw(n) = O\left(\frac{2^n (\ln n)^2}{n}\right) = o(2^n)$.

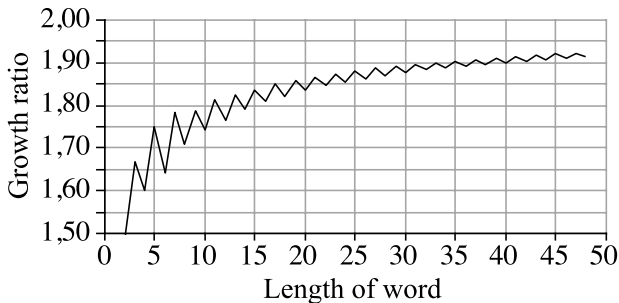
Enumerating p.n. words

Thanks to the new generating algorithm, we can now count p.n. words up to $n = 50$, and also special subsets, e.g. fixed-density, fixed-prefix, fixed-suffix, . . . And make intelligent guesses.

Enumerating p.n. words

OEIS sequence no. A194850

| n | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
|----------|---|---|---|---|----|----|----|----|-----|-----|-----|-----|------|------|------|------|
| $pnw(n)$ | 2 | 3 | 5 | 8 | 14 | 23 | 41 | 70 | 125 | 218 | 395 | 697 | 1273 | 2279 | 4185 | 7568 |



Fixed-density p.n. words

For a word w , its **density** d is the number of 1s in w ,
 $pnw(n, d) =$ no. of p.n. words of length n and density d .

Let $f_d(x) = \sum_{n \geq 0} pnw(n, d)x^n$, the generating function of $pnw(n, d)$.

| gen.func. | coefficients | closed form |
|---------------------------------------|--------------------------------|--|
| $f_0(x) = \frac{1}{1-x}$ | 1, 1, 1, 1, 1, 1, 1, 1, 1, ... | 1 |
| $f_1(x) = \frac{x}{1-x}$ | 0, 1, 1, 1, 1, 1, 1, 1, 1, ... | $[n > 0]$ |
| $f_2(x) = \frac{x^2}{(1-x)^2}$ | 0, 0, 1, 2, 3, 4, 5, 6, ... | $(n-1)$ |
| $f_3(x) = \frac{x^3}{(1-x^2)(1-x)^2}$ | 0, 0, 0, 1, 2, 4, 6, 9, ... | $pnw(2n) = n(n-1),$ $pnw(2n+1) = (n-1)^2$ |

... and more (but for larger d becomes less and less manageable) ...

prefix **normal** not to be confused with **abnormal**,



prefix normal not to be confused with abnormal, AB normal,



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Mel Brooks: Young Frankenstein (1974)

THANK YOU!