



A suffix-array based algorithm for expression clustering

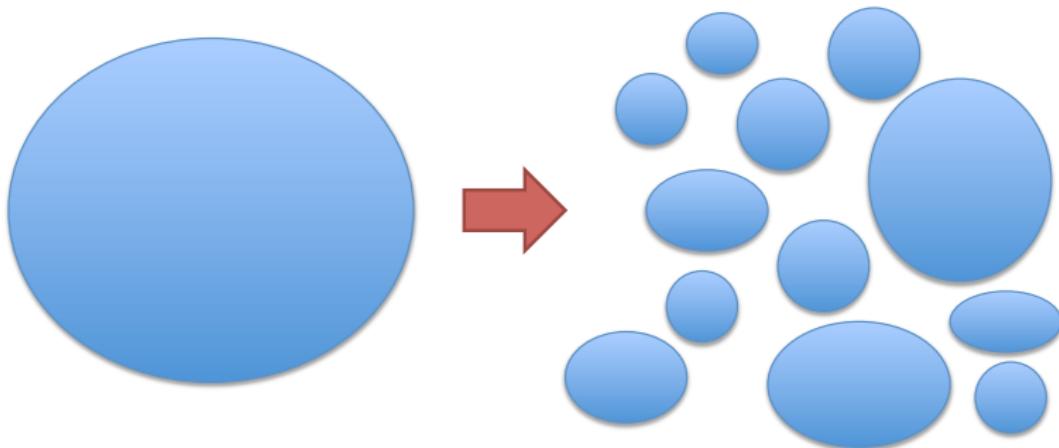
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University of the Witwatersrand, Johannesburg (South Africa)

Zsuzsanna Lipták

University of Salerno/University of Verona (Italy)

What is the problem?



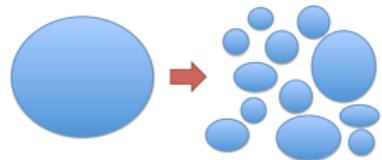
Given:

Set of n sequences (10^5 or more) of length m

Wanted:

clusters s.t. each cluster corresponds to a gene

Using pairwise comparison



$n = \# \text{ seq's}, m = \text{ave. length}$

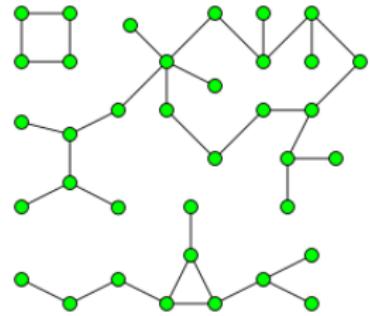
Single linkage

s, t similar $\Rightarrow s, t$ in same cluster

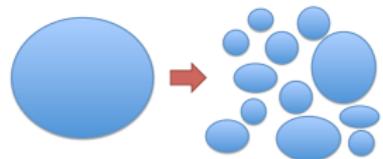
As a graph problem

Find connected components.

Remember: Edges have to be computed!



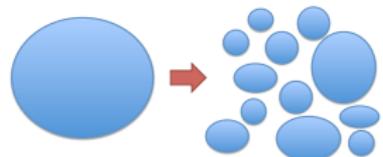
Using pairwise comparison



$n = \#$ seq's, $m = \text{ave. length}$

- One comparison (e.g. edit dist.) costs m^2 $\rightsquigarrow O(n^2 m^2)$ time

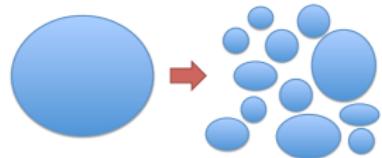
Using pairwise comparison



$n = \#$ seq's, $m = \text{ave. length}$

- One comparison (e.g. edit dist.) costs $m^2 \rightsquigarrow O(n^2 m^2)$ time
Most pairs are not similar: use filters e.g. $O(m)$ time
if x, y pass filter then compute $d(x, y)$

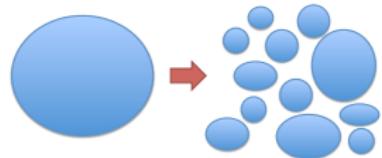
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- If one filter test costs $m \rightsquigarrow O(n^2 m)$ time + sg.

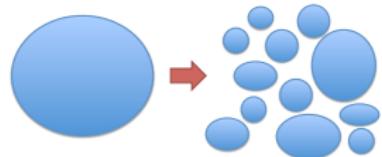
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- If one filter test costs $m \rightsquigarrow O(n^2 m)$ time + sg.
- In general:
If one filter test costs $f(m) \rightsquigarrow O(n^2 f(m))$ time + sg.

Using pairwise comparison



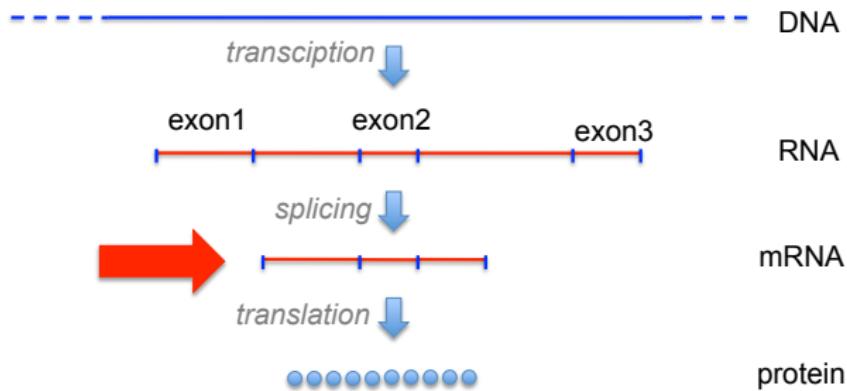
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- In general:
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Here: We get rid of the n^2 factor!

Clustering expression data

Expression data (formerly ESTs)



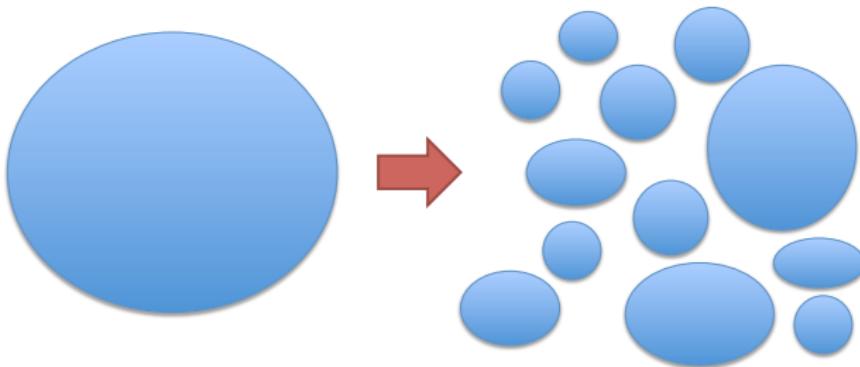
Transcriptome/Expression data

(ESTs = expressed sequence tags)

Partial cDNA copies of mRNAs (lab process)

both Sanger-style (ESTs) and NGS

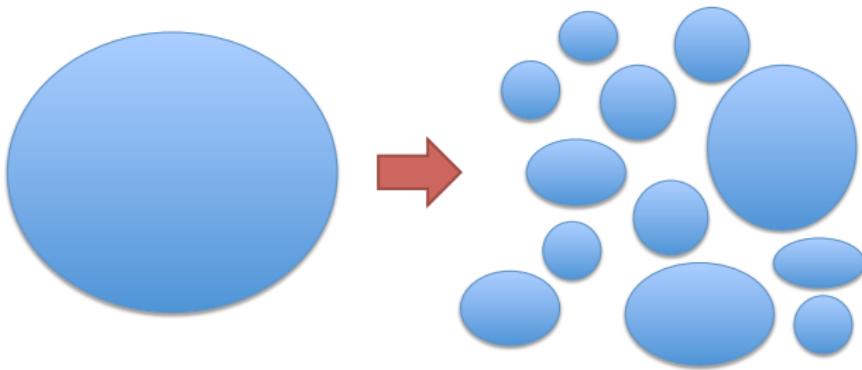
Problem statement



Desideratum

Given a set of sequences S , find a partition of S s.t. if s, t are products of the same gene, then s, t in same cluster.

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Given a set of sequences S , find a partition of S s.t. if s, t are products of the same gene, then s, t in same cluster.

Formal Problem

Given a set of sequences S , find a partition of S s.t. if s, t are **similar**, then s, t in same cluster.

Expression clustering

- expression studies (e.g. diseased vs. healthy),
- gene discovery,
- SNP detection,
- discovery of products of **alternative splicing**,
- estimating no. of genes

We are not doing:

- consensus building
- alignment to genome

String similarity

Example

$s = \text{CAAGACAA}$, $t = \text{CAGAGCAC}$

- alignment / edit distance (N/W 1970, S/W 1981)
- q -gram-distance: $\sum_{|w|=q} |\text{freq}(w,s) - \text{freq}(w,t)|$ (Ukkonen 1992)
- d^2 : $\sum_{|w|=q} (\text{freq}(w,s) - \text{freq}(w,t))^2$ (Torney et al. 1990)
- others (fingerprints, information theory based, ...)

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CAAGA–CAA
CA–GAGCAC

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CAAGA–CAA
CA–GAGCAC

AA	2	0
AC	1	1
AG	1	2
CA	2	2
GA	1	1
GC	0	1

String similarity

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$s = \text{CAAGACAA}$, $t = \text{CAGAGCAC}$

- alignment / edit distance = 3 (N/W 1970, S/W 1981)
- q -gram-distance: $\sum_{|w|=q} |\text{freq}(w,s) - \text{freq}(w,t)| = 4$ (Ukkonen 1992)
- d^2 : $\sum_{|w|=q} (\text{freq}(w,s) - \text{freq}(w,t))^2$ (Torney et al. 1990)
- others (fingerprints, information theory based, ...)

CAAGA–CAA
CA–GAGCAC

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AG	1	2	1
CA	2	2	0
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- d^2 : $\sum_{|w|=q} (\text{freq}(w,s) - \text{freq}(w,t))^2 = 6$ (Torney et al. 1990)
- others (fingerprints, information theory based, ...)

CAAGA–CAA
CA–GAGCAC

AA	2	0	2	4
AC	1	1	0	0
AG	1	2	1	1
CA	2	2	0	0
GA	1	1	0	0
GC	0	1	1	1

4 6

Similar windows

>1

tttattattatgggattaggcttctctcgcgattctttacgtt
ggttatatctttatattacttttttttattgtataatttagcttccg
tcctctaggtcccttctcta

>2

ccccctctgttggttatctttatattacttttttttattgtataaa
cttcttattattctccgcctctctttatcttccggattctcgagg
ggatatatattatgcggcgccctttatatctctatattctctc

a

Similar windows

>1

tttattattatgggattaggcttctctcgcgattctttacGTT
CGTTATATCTTATATTACTTTCTTATTGTATAAttagcttccg
tcctctaggtcccttctcta

>2

cccctctctGTTGGTTATATCTTATATTACTTCTTATTGTATAA
cttcttattattctccgctctctcttatcttccggaattctcgagg
ggatatatatattatgcggcgccctcttatatctctatattctctc

a

Our string similarity measure

We use d^2 , minimized over all pairs of windows of given size.

$$d(s, t) = \min_{s' \sqsubseteq s, t' \sqsubseteq t, |s|=|t|=\mu} \sum_{|w|=q} (\text{freq}(w, s') - \text{freq}(w, t'))^2$$

Parameters:

wordsize q ($= 6, 8$), windowsize, threshold (s, t similar iff $d(s, t) < \theta$)

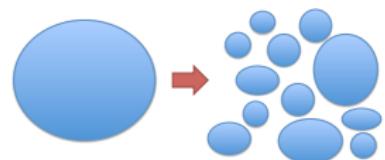
Computation time: $O(m^2)$

Speedup of previous tool

wcd + KABOOM! = wcd-express

The KABOOM-filter

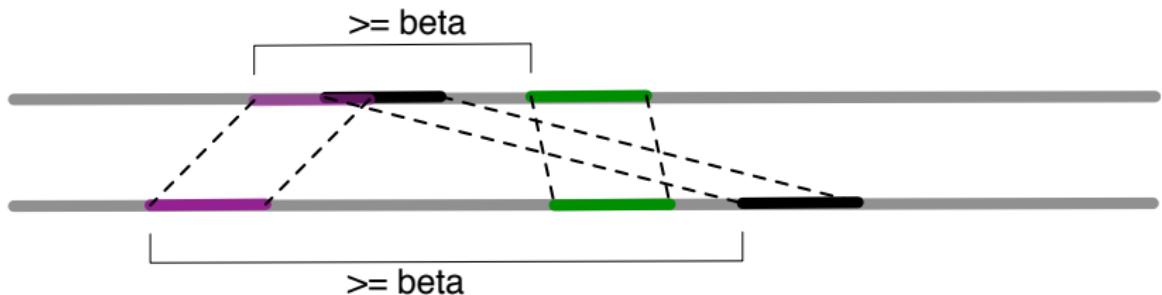
Filters for d^2



$n = \#$ seq's, $m = \text{ave. length}$

- Recall: we want to avoid $O(n^2 m^2)$ runtime
- Filters with guarantee that need pairwise comparison only reduce m^2 factor
- One long exact match can be implemented with suffix array (linear in nm) – but too fragile!
- KABOOM-filter combines good running time with good filtering qualities

The KABOOM-filter

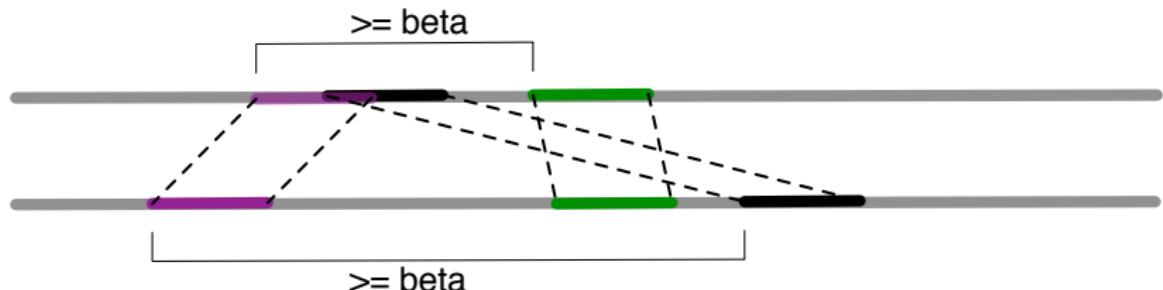


Definition

s, t are (k, α, β) -similar if they have at least α many k -words in common, at least β apart.

- k -words are counted with multiplicities (a.k.a. D^2)
- distance β between leftmost and rightmost occurrence, in both strings
- order of matches not important

The KABOOM-filter



Definition

s, t are (k, α, β) -similar if they have at least α many k -words in common, at least β apart.

- heuristic filter for d^2
- very good in practice: produces practically no FNs and reduces no. of d^2 comparisons dramatically
- $k > q$ (typically k between 12 and 20)

Implementation of the KABOOM-filter

```
for every sequence  $s_i$ 
    for every  $k$ -word  $w$  in  $s_i$ 
        find all  $s_j$  containing  $w$ 
         $M = M \cup s_j$                                 ( $M$ : candidate sequences for  $s_i$ )
        update variables for  $s_j$ 
    for every  $s_j$  in  $M$ 
        check whether  $s_i, s_j$  fulfill  $(k, \alpha, \beta)$ -criterion
        if yes, compute  $d(s_i, s_j)$ 
```

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```

We need

- suffix array of concatenated string
 - inverse of suffix array

Implementation of the KABOOM-filter

T27882	g609980 T27882 CLONE_LIB: Human Eye.	2009	AGCTTATTTC....CCAA	661
	TGCTTCAAACGGGTCTGGAATTTCTCAAAGGTATGGTGAAGAATAT	2010	AGCTTATTTC....CCCA	555
	ATAGTCCAGGGACGAATCAAATGGTAGAGACTCATGAAGATGCAC	2011	AGCTTATTTC....CCCC	7
	GCTGGGTGACATGGACATCTCACCGCTGGGACTTCTCGAAGACCA	2012	AGCTTATTTC....CCCC	1
	TACCCAGAAATCTGACACTCAAGACATACCTG	2013	AGCTTATTTC....CCCC	130
T27883	g609981 T27883 CLONE_LIB: Human Eye.	2014	AGCTTATTTC....CCCC	18
	CTGGGCAAATCTAACGCTGCCAGGGCAGTTAGGTTGACCGCTGTNTGTT	2015	AGCTTATTTC....CCCC	400
	CGAGGGCAGGGGGGAGTCATCACCTCTGGCTGGACCGGGCTGCTG	2016	AGCTTATTTC....CCCT	
	TCTGCTGAACTCTGGGAGTCTGCACTAACCCGGGAGTGGTGTCTCCAIG			
	AGAGCTTGGACCAAGATITGATGATCTACCTGAGCTGGGCCATGTTGTTCTG			
			
T27884	g609982 T27884 CLONE_LIB: Human Eye.			
	CCAGCAGACAGAGGACTCTATTAAGGAAGGTCTGTGCGCTGACCTAC			
	AGAGAAGATGCTACTCATCTATGGTACCCAGCTGGGCCATGTTGTTCTG			
	GGG			
T27885	g609983 T27885 CLONE_LIB: Human Eye.			
	CCCAAACCGATAAAGTGCATGACGAAACAATAGGACTTCCCAGGGNTGAA			
	TCCCGGCCAGCCAGGTGACCCCAAGGTCTGGATGNTCTGGCTGTTCTTC			
T27886	g609984 T27886 CLONE_LIB: Human Eye.			
	ATCGCGCCATCATGGCTGCATGCATGCTCCGGAAAGGGCTGTCCAGT			
	ATTCTNANTAGAGGCGGATTTACCGCTTGGCTTGATATAAGACC			
	TCTCCCTCCAAATGGAAATATGGAAATCATCTACAGCCTCTGCCCTGGT			
	TTTGTCTGTACTCAAGCAATAAATGGTTGTTAAC			
T27887	g609985 T27887 CLONE_LIB: Human Eye.			
	GGCACAGTCAGCTTCTGGGGCCGGTGTATGTTGGCTCTGGGGCGGG			
	CTGGNATTCACGGATTCGGCTGCCACCCAGATCCANCTGNTCCACTTCAC			
	TCTGGTGTGATTAATGAGGGCAGGGG			
T27888	g609986 T27888 CLONE_LIB: Human Eye.			
	ACACCTGCCATGTCAGCATGAGGTTGGCCAAAGCCCTCACCTGAGATG			
T27890	g609988 T27890 CLONE_LIB: Human Eye.			
	CCGCAGNTGCAAGAAAANTTGGNGGTGATGAGGGACATCAGTGANICA			
	TGGTATGGGGACAAAAAAGRCTTAGAAGATGGAGATCAACCGAGATGCT			
	TGCTCTCAAAAGACTCTTGGAACACAGTACACCGATCAGCAG			
	TTAATG			

Modified suffix array sa_k

aaa@aacgggt@gttaaagt@tcgggt@gttat@cgg@acgggt@

i	sa	Text from $sa[i]$	sa_k
15	16	agt@tcgggt@gttat@cgg@acgggt@	16
16	29	at@cgg@acgggt@	29
17	32	cgg@acgggt@	6
18	37	cgggt@	21
19	6	cgggt@gttaaagt@t...gt@	32
20	21	cgggt@gttat@cgg@acgggt@	37
21	34	g@acgggt@	34
22	33	gg@acgggt@	33
23	38	ggt@	7
24	7	ggt@gttaaagt@t...t@	22
25	22	ggt@gttat@cgg@acgggt@	38
26	39	gt@	8

Modified suffix array sa_k

aaa@aa**cgg**t@gtt@aaagt@**cgg**t@gtt@at@**cgg**@a**cgg**t@

i	sa	Text from $sa[i]$	sa_k
15	16	agt@tcggt@gttat@cgg@acggt@	16
16	29	at@cgg@acggt@	29
17	32	cgg@acggt@	6
18	37	cggt@	21
19	6	cggt@gtt@aaagt@t...gt@	32
20	21	cggt@gtt@at@cgg@acggt@	37
21	34	g@acggt@	34
22	33	gg@acggt@	33
23	38	ggt@	7
24	7	ggt@gtt@aaagt@t...t@	22
25	22	ggt@gtt@at@cgg@acggt@	38
26	39	gt@	8

Analysis

Time

- $\sum_{|w|=k} \text{freq}(w)^2$, which is between nm and n^2m^2
- in practice: very fast
- WE work efficiency: $1 - (\text{TP} + \text{FP}) / (\text{all pairs})$ above 99.5%

Space

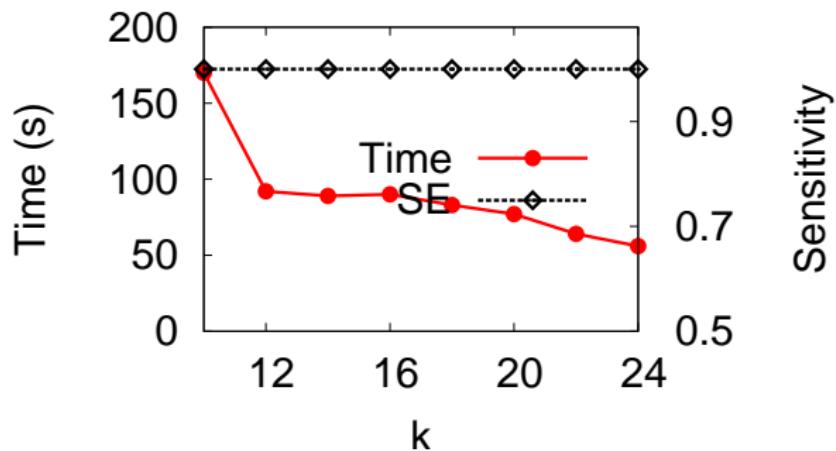
- text nm (nt)
- k -depth suffix array: $nm \log n$ ($\times 2$ for RC)
- inverse: $nm \log n$

Experimental results

Speedup w.r.t. wcd

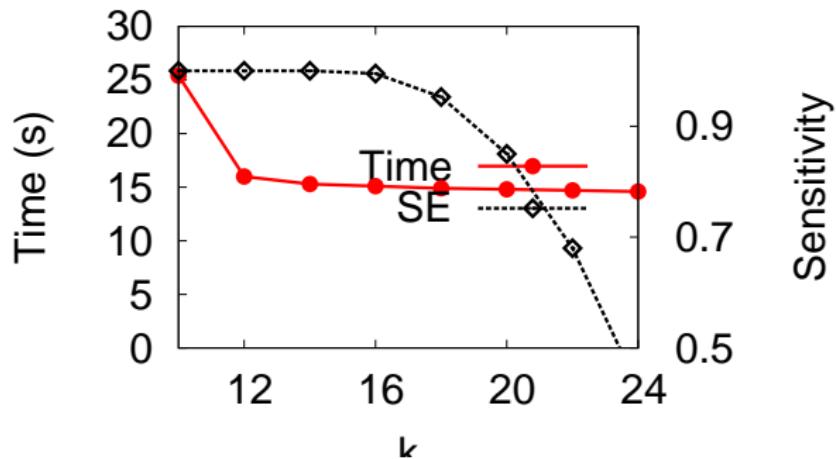
Data set	# seqs (K)	size Mb	γE	wcd-exp (s)	wcd (s)	Speed-up factor
A076941	77	32	17	100	578	5.7
A208	484	208	102	1240	23983	19.3
C10	126	56	355	511	4512	8.8
chlamy	190	100	139	1000	5989	6.0
Drosophila	25	86	68	184	1542	8.4
ecHuman	17	11	171	135	496	3.7
pubcot	30	17	34	65	222	3.4
ricinus	58	40	162	840	1518	1.8
xen	233	137	63	855	9298	10.8

Sensitivity to choice of k



Filtered cotton set, $\alpha = 3, \beta = 32, H = 48$.

Sensitivity to choice of k



Synthetic 454 data set, $\alpha = 3$, $\beta = 16$, $H = 32$.

Comparison with other tools

Issues

- choice of parameters
- pre-/postprocessing (here: filtering)
- quality: what's the truth?
- synthetic vs. real data

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Citroën/Ferrari comparison problem



FASTER THAN A FERRARI.

Travelling flat out at 71.5mph the Citroen 2CV will easily overtake the Ferrari Mondial travelling at 65mph.



AS MANY WHEELS AS A ROLLS ROYCE.

The £23,440 Rolls-Royce Silver Spirit. How many wheels? Four. The £2,584 Citroen 2CV. How many wheels? Exactly the same.



MORE ROOM THAN A PORSCHE.

With a possible 30cu. ft. boot space there's no need for one of those plastic luggage racks on our little run-about.



THE £2,674 CITROËN 2CV.

All you'll ever need in a car.

PRICE TELLED BY YOUR CITROËN DEALER. 1984. CAR SHOWN WITH 1.1L PETROL ENGINE. ALSO AVAILABLE WITH 1.3L DIESEL ENGINE. FOR NEAREST DEALER
TOLL FREE 1-800-227-2277. CITROËN, THE CITROËN LOGO AND CITROËN 2CV ARE TRADEMARKS OWNED BY CITROËN S.A. ©1984 CITROËN S.A.

PUNCH September 19 1984

Quality and runtime comparison

	Sensitivity		Jaccard Index		Time (s)	
	wcdx	PEACE	wcdx	PEACE	wcdx	PEACE
A076941	0.932	0.930	0.473	0.477	100	951
chlamy	0.949	0.949	0.513	0.513	907	8823
ecHuman	0.996	0.998	0.707	0.630	50	147
metasim454	0.793	0.714	0.765	0.689	16	66
metasimIllum	0.444	0.368	0.398	0.364	19	1975

$$\text{SE} = \text{TP}/(\text{TP}+\text{FN}), \text{JI} = \text{TP}/(\text{TP}+\text{FP}+\text{FN}), \text{PPV} = \text{TP}/(\text{TP}+\text{FP})$$

PEACE: *NAR*, 2010

Quality of clustering using KABOOM-heuristic

“We can drive our car at least as fast as others have driven their car.”

- ESTate
- PaCE
- PEACE
- TIGR

Quality of clustering using KABOOM-heuristic

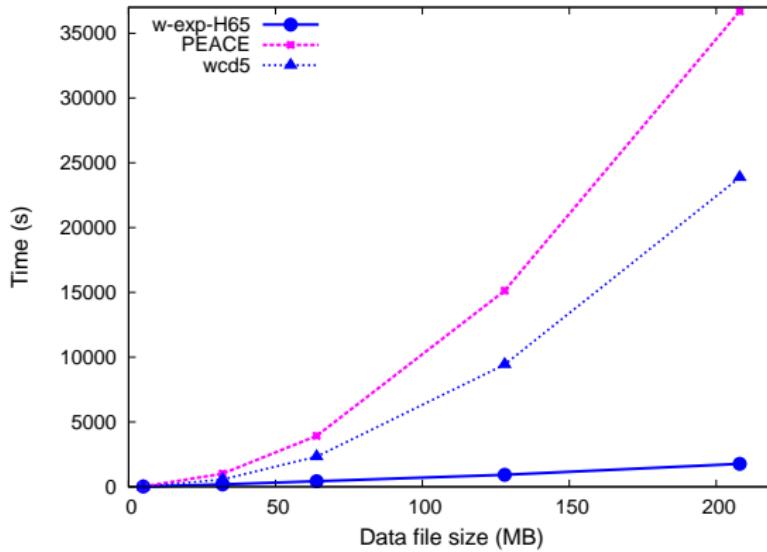
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- ESTate
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- TIGR

Performance matters.

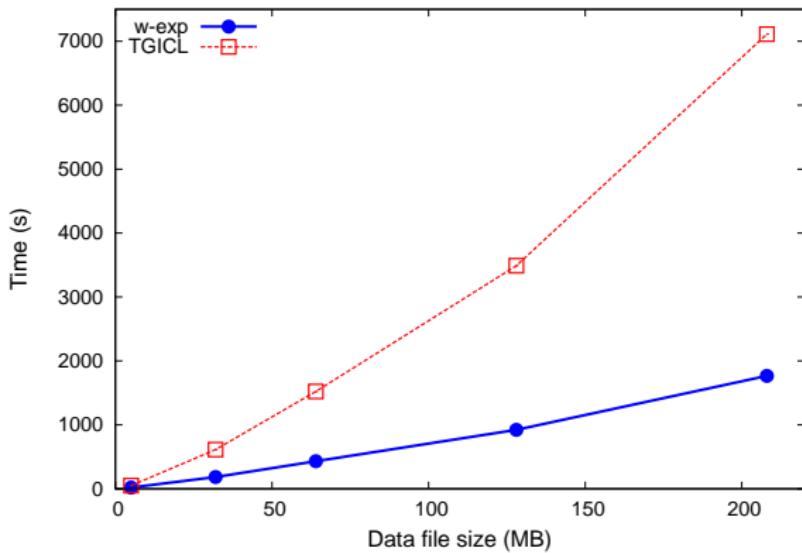
Runtime comparison

Arabidopsis – wcd-express versus PEACE



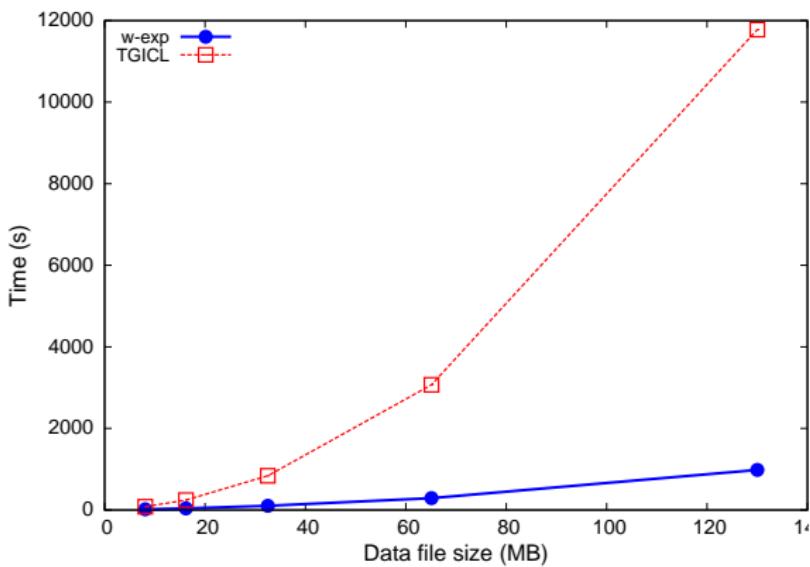
Runtime comparison

Arabidopsis – wcd-express versus TGICL



Runtime comparison

Human 454 ESTs – wcd-express versus TIGR



Conclusion

- wcd + KABOOM! = wcd-express
- Much faster and at least as good as other tools.
- Memory: very heavy memory usage (suffix array + inverse)

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Ongoing/future work

- Memory needs to be reduced!
- Integrate wcd-express into workflow (filtering/masking, assembly)
- GUI?

Paper to appear in Bioinformatics.



it starts with a
playground.

Thank you.

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