

# An Introduction to Model Theory for Representation Theory

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## Definition

A (first-order) **language**  $\mathcal{L}$  consists of

- ▶ 3 mutually disjoint sets:  $\mathcal{R}$  the set of **relation symbols**,  $\mathcal{F}$  the set of **function symbols** and  $\mathcal{C}$  the set of **constant symbols**; and
- ▶ an **arity function**  $\lambda : \mathcal{R} \cup \mathcal{F} \rightarrow \mathbb{N}$ .

For any  $Q \in \mathcal{R} \cup \mathcal{F}$ , we will refer to  $\lambda(Q)$  as the **arity** of  $Q$ .

## Definition

An  $\mathcal{L}$ -**structure**  $\mathcal{A}$  is a non-empty set  $A$  called the **domain** of  $\mathcal{A}$  together with

- (i) a subset  $R^{\mathcal{A}}$  of  $A^{\lambda(R)}$  for each  $R \in \mathcal{R}$ ;
- (ii) a function  $F^{\mathcal{A}}$  from  $A^{\lambda(F)} \rightarrow A$  for each  $F \in \mathcal{F}$ ; and
- (iii) an element  $c^{\mathcal{A}} \in A$  for each  $c \in \mathcal{C}$ .

For  $Q \in \mathcal{R} \cup \mathcal{F} \cup \mathcal{C}$ , we call  $Q^{\mathcal{A}}$  the **interpretation** in  $\mathcal{A}$ .

## Example

Let  $\mathcal{L} := \{\circ\}$  where  $\circ$  is a binary function symbol. An  $\mathcal{L}$ -structure is a magma.

## More examples of languages

- The language of abelian groups is  $\mathcal{L}_{Ab} := \{0, +, -\}$  where  $0$  is a constant symbol,  $+$  is a binary function symbol and  $-$  is unary function symbol.
- The language of rings is  $\mathcal{L}_{rings} := \{0, 1, +, -, \cdot\}$  where  $0$  is a constant symbol,  $+$  is a binary function symbol and  $+$  and  $\cdot$  are binary function symbols and  $-$  is a unary function symbol.
- For  $R$  a ring, the language of  $R$ -modules is  $\mathcal{L}_R := \{0, +, (\cdot r)_{r \in R}\}$  where  $0$  is a constant symbol,  $+$  is a binary function symbol and for each  $r \in R$ ,  $\cdot r$  is a unary function symbol.
- The language of ordered sets is  $\mathcal{L}_{\leq} := \{\leq\}$  where  $\leq$  is a binary relation symbol.

## Definition

The **alphabet** of a language  $\mathcal{L}$  is the relation, functions and constant symbols of  $\mathcal{L}$  together with a set of logical symbols which are part of every language consisting of:

Connectives:  $\{\rightarrow, \wedge, \vee, \neg\}$

Quantifiers:  $\forall$  and  $\exists$

The equality symbols  $=$

Brackets “)” and “(”

Comma: “,”

A set of variables denoted  $V_{bl} := \{v_i \mid i \in \mathbb{N}\} \cup \{u, v, w, x, y, z\}$

## Examples of $\mathcal{L}$ -formulae:

The  $\mathcal{L}_{rings}$ -formula

$$(\forall v_2 \ v_1 \cdot v_2 = v_2 \cdot v_1)$$

defines the centre of a ring.

The  $\mathcal{L}_{rings}$ -formula

$$(\forall v_1 (v_1 = 0 \vee (\exists v_2 \ v_1 \cdot v_2 = 1)))$$

expresses that every non-zero element is invertible.

Define  $\text{tm}_0(\mathcal{L})$  to be the set  $\text{Vbl} \cup \mathcal{C}$ . For all  $k \in \mathbb{N}$ , let  $\text{tm}_{k+1}(\mathcal{L})$  be  $\text{tm}_k(\mathcal{L}) \cup \{F(t_1, t_2, \dots, t_n) \mid F \in \mathcal{F}, \lambda(F) = n \text{ and } t_1, \dots, t_n \in \text{tm}_k(\mathcal{L})\}$ .

We define the set of  $\mathcal{L}$ -**terms** to be

$$\text{tm}(\mathcal{L}) := \bigcup_{k \in \mathbb{N}_0} \text{tm}_k(\mathcal{L}).$$

Let  $\text{Fml}_0(\mathcal{L})$  be the set of strings in the alphabet of  $\mathcal{L}$  of the form

$$t_1 = t_2 \text{ or } R(t_1, \dots, t_n)$$

where  $t_1, \dots, t_n$  are  $\mathcal{L}$ -terms and  $R \in \mathcal{R}$  has arity  $n$ .

For each  $k \in \mathbb{N}_0$ , let

$$\begin{aligned} \text{Fml}_{k+1}(\mathcal{L}) := & \text{Fml}_k(\mathcal{L}) \cup \{(\varphi \rightarrow \psi), (\varphi \wedge \psi), (\varphi \vee \psi), \neg\varphi \mid \varphi, \psi \in \text{Fml}_k(\mathcal{L})\} \\ & \cup \{(\forall x\varphi), (\exists x\varphi) \mid \varphi \in \text{Fml}_k(\mathcal{L}) \text{ and } x \in \text{Vbl}\}. \end{aligned}$$

We define the set of  $\mathcal{L}$ -**formulae** to be

$$\text{Fml}(\mathcal{L}) := \bigcup_{i \in \mathbb{N}_0} \text{Fml}_i(\mathcal{L}).$$

## Interpreting formulae in $\mathcal{L}$ -structures

Let  $\theta$  be an  $\mathcal{L}$ -formula and

$$(\exists v_i \text{ \underline{\hspace{1cm}} }) \quad \text{or} \quad (\forall v_i \text{ \underline{\hspace{1cm}} })$$

a substring of  $\theta$ .

**Free variables:** An instance of a variable is **free** if it is not a bound instance or a quantifier instance. The **free variables** of a formula  $\theta$  are those variable which occur as free instances.

To indicate that an  $\mathcal{L}$ -formula  $\theta$  has free variables contained in a set  $\{x_1, \dots, x_n\}$  we write  $\theta(x_1, \dots, x_n)$ .

Let  $M$  be an  $\mathcal{L}$ -structure,  $\theta(x_1, \dots, x_n)$  be an  $\mathcal{L}$ -formula and  $(m_1, \dots, m_n) \in M$ . We say that  $\theta(m_1, \dots, m_n)$  holds in  $M$  if the expression obtained by replacing every free instance of  $x_i$  in  $\theta$  by  $m_i$  is true in  $M$ .

A **sentence** is a formula without free variables.

The  $\mathcal{L} = \{\circ\}$ -formula

$$(\forall x(\forall y(\forall z(x \circ y) \circ z = x \circ (y \circ z))))$$

is a sentence.

## Interpreting formulae in $\mathcal{L}$ -structures

Let  $\theta$  be an  $\mathcal{L}$ -formula and

$$(\exists v_i \text{ ---}) \quad \text{or} \quad (\forall v_i \text{ ---})$$

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A **sentence** is a formula without free variables.

Let  $\theta(v_1)$  be the  $\mathcal{L}_{rings}$ -formula

$$(\forall v_2 \ v_1 \cdot v_2 = v_2 \cdot v_1)$$

is not a sentence.

If  $\Sigma$  is a set of sentences then we say an  $\mathcal{L}$ -structure  $\mathcal{M}$  is a **model** of  $\Sigma$  if all sentences in  $\Sigma$  hold in  $\mathcal{M}$ .

If  $X$  is a class of  $\mathcal{L}$ -structures then the **theory of  $X$** , written  $\text{Th}(X)$ , is the set of all  $\mathcal{L}$ -sentences which hold in all members of  $X$ .

A class  $X$  of  $\mathcal{L}$ -structures is **axiomatisable** if there is a set of sentences  $\Sigma$  such that the members of  $X$  are exactly the models of  $\Sigma$ . **Examples**

We say  $\mathcal{L}$ -structures  $\mathcal{M}$  and  $\mathcal{N}$  are **elementary equivalent**, and write  $\mathcal{M} \equiv \mathcal{N}$ , if  $\text{Th}(\mathcal{M}) = \text{Th}(\mathcal{N})$ .

Let  $\mathcal{M}$  be an  $\mathcal{L}$ -structure with domain  $M$ . A subset  $N \subseteq M$  is an **elementary** if for all formulas  $\theta(x_1, \dots, x_n)$  and  $(a_1, \dots, a_n) \in N$  of elements of  $N$ ,  $\theta(a_1, \dots, a_n)$  holds in  $N$  if and only if  $\theta(a_1, \dots, a_n)$  holds in  $M$ .



# The First Theorem of Model Theory

## The Compactness Theorem

Let  $\mathcal{L}$  be a language.

- (i) A set  $\Sigma$  of  $\mathcal{L}$ -sentences has a model if and only if every finite subset of  $\Sigma$  has a model.
- (ii) Let  $\Sigma$  be a set of formulas with free variables  $(v_1, v_2, \dots)$ .

If for every finite subset  $\Sigma'$  of  $\Sigma$ , there exists an  $\mathcal{L}$ -structure  $M$  and a tuple of elements  $\bar{m}$  such that  $\theta(\bar{m})$  holds in  $M$

then there exists an  $\mathcal{L}$ -structure  $M$  and a tuple of elements  $\bar{m}$  such that  $\theta(\bar{m})$  holds in  $M$  for all  $\theta \in \Sigma$ .

# Model Theory of Modules

Let  $R$  be a ring and  $\mathcal{L}_R := \{0, +, (\cdot r)_{r \in R}\}$  the language of  $R$ -modules.

Every  $\mathcal{L}_R$ -term is equivalent (relative to  $\text{Th}(\text{Mod-}R)$ ) to one of the form

$$\sum_{i=1}^n x_i \cdot r_i$$

where each  $x_i$  is a variable and each  $r_i \in R$  for  $1 \leq i \leq n$ .

Every atomic formulae is equivalent (relative to  $\text{Th}(\text{Mod-}R)$ ) to one of the form

$$\sum_{i=1}^n x_i \cdot r_i = 0$$

(or  $0 = 0$ ) where  $x_1, \dots, x_n$  are variables.

A (right) **pp- $n$ -formula** (over  $R$ ) is a formula  $\varphi(\bar{x})$  of the form

$$\exists y_1, \dots, y_m \bigwedge_{i=1}^l \sum_{j=1}^n x_j r_{ij} + \sum_{k=1}^m y_k s_{ik} = 0$$

where  $r_{ij}, s_{ik} \in R$  and  $\bar{x} = (x_1, \dots, x_n)$ .

For  $M \in \text{Mod-}R$ , we write  $\varphi(M)$  for the solution set of  $\varphi$  in  $M$ .

- If  $m_1, m_2 \in \varphi(M)$  then  $m_1 + m_2 \in \varphi(M)$ .
- If  $f : M \rightarrow L \in \text{Mod-}R$  and  $\bar{m} \in \varphi(M)$  then  $f(\bar{m}) \in \varphi(L)$ .
- Let  $N_i \in \text{Mod-}R$  for  $i \in I$ . Then

$$\varphi(\oplus_i N_i) = \oplus_i \varphi(N_i).$$

A (right) **pp- $n$ -formula** (over  $R$ ) is a formula  $\varphi(\bar{x})$  of the form

$$\exists y_1, \dots, y_m \bigwedge_{i=1}^l \sum_{j=1}^n x_j r_{ij} + \sum_{k=1}^m y_k s_{ik} = 0$$

where  $r_{ij}, s_{ik} \in R$  and  $\bar{x} = (x_1, \dots, x_n)$ .

We write  $\text{pp}_R^n$  for the set of (right) pp- $n$ -formulae over  $R$  where we identify pp- $n$ -formulae  $\varphi, \psi$  if  $\varphi(M) = \psi(M)$  for all  $M \in \text{Mod-}R$ .

$\text{pp}_R^n$  is a bounded modular lattice when equipped with the order defined by

$$\psi \leq \varphi \text{ if and only if } \psi(M) \subseteq \varphi(M) \text{ for all } M \in \text{Mod-}R.$$

We write  $\varphi + \psi$  for the join (l.u.b) and  $\varphi \wedge \psi$  for the meet (g.l.b) in  $\text{pp}_R^n$ .  
For all  $M \in \text{Mod-}R$ ,

$$(\varphi + \psi)(M) = \varphi(M) + \psi(M) \text{ and } (\varphi \wedge \psi)(M) = \varphi(M) \cap \psi(M).$$

**Modular:**  $a \leq b$  implies  $a + (z \wedge b) = (a + z) \wedge b$ .



Let  $\varphi, \psi \in \text{pp}_R^n$  with  $\psi \leq \varphi$  and let  $b \in \mathbb{N}$ . We write

$$|\varphi/\psi| \geq b$$

for the  $\mathcal{L}_R$ -sentence which expresses in all  $R$ -modules  $M$  that

$$|\varphi(M)/\psi(M)| \geq b.$$

Suppose  $n = 1$ . Then we may take  $|\varphi/\psi| \geq b$  to be

$$\exists z_1, \dots, z_b \bigwedge_{i=1}^b \varphi(z_i) \wedge \bigwedge_{i < j} \neg \psi(z_i - z_j).$$

## The Baur-Monk Theorem

Every formula in the language of  $R$ -modules is equivalent to a boolean combination of pp-formulae and sentences of the form

$$|\varphi/\psi| \geq b$$

where  $b \in \mathbb{N}$  and  $\varphi, \psi$  are pp-1-formulae such that  $\psi \leq \varphi$ .

## Corollary

Let  $M, N \in \text{Mod-}R$ . Then

$$M \equiv N \quad \text{if and only if} \quad |\varphi(M)/\psi(M)| = |\varphi(N)/\psi(N)|,$$

when either is finite, for all  $\varphi \geq \psi \in \text{pp}_R^1$ .

## Examples

- Let  $N_i$  be a collection of  $R$ -modules for  $i \in I$ , then

$$\bigoplus_{i \in I} N_i \equiv \prod_{i \in I} N_i.$$

- If  $R$  is an algebra over an infinite field  $k$  then for all  $M \in \text{Mod-}R$ ,  $M^2 \equiv M$ .
- As  $\mathbb{Z}$ -modules,  $\mathbb{Z} \oplus \mathbb{Q} \equiv \mathbb{Z}$ .

## Corollary

A submodule  $L \subseteq M$  is elementary if and only if  $L \equiv M$  and, for all pp-formulae  $\varphi$ ,  $\varphi(L) = \varphi(M) \cap L$ .

# Purity

## Definition

An embedding  $f : M \hookrightarrow N$  is **pure** if for all pp-1-formulae  $\varphi$ ,

$$\varphi(N) \cap f(M) = f(\varphi(M)).$$

An  $R$ -module  $M$  is **pure-injective** if every pure-embedding  $M \hookrightarrow N$  splits.

## Definition

An  $R$ -module  $M$  is **algebraically compact** if any system of (inhomogeneous) linear equations over  $R$ , in arbitrary many variables, which is finitely solvable in  $M$ , has a solution in  $M$ .

Equivalently, an  $R$ -module  $M$  is algebraically compact if for any  $n \in \mathbb{N}$ , the collection of sets of the form  $\bar{a} + \varphi(M)$  where  $\varphi \in \text{pp}_R^n$  and  $\bar{a}$  is an  $n$ -tuple in  $M$ , has the finite intersection property.

## Theorem

*An  $R$ -module is algebraically compact if and only if it is pure-injective.*

# Modules up to Elementary Equivalence

Write  $\text{pinj}_R$  for the set of indecomposable pure-injective  $R$ -modules.

## Theorem (Ziegler)

*For every  $R$ -module  $M$ , there exists  $N_i \in \text{pinj}_R$  such that  $M$  is elementary equivalent to  $\bigoplus_{i \in I} N_i$ .*

## Theorem

*Let  $N, M \in \text{pinj}_R$ . Then  $N \equiv M$  if and only if for all  $\varphi \geq \psi \in \text{pp}_R^1$*

$$|\varphi(N)/\psi(N)| > 1 \Leftrightarrow |\varphi(M)/\psi(M)| > 1.$$



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## Reminders

A (right) **pp- $n$ -formula** (over  $R$ ) is a formula  $\varphi(\bar{x})$  of the form

$$\exists y_1, \dots, y_m \bigwedge_{i=1}^l \sum_{j=1}^n x_j r_{ij} + \sum_{k=1}^m y_k s_{ik} = 0$$

where  $r_{ij}, s_{ik} \in R$  and  $\bar{x} = (x_1, \dots, x_n)$ .

We write  $\text{pp}_R^n$  for the set of (right) pp- $n$ -formulae over  $R$ . This set is a lattice when ordered by  $\varphi \geq \psi$  if and only if  $\varphi(M) \supseteq \psi(M)$  for all  $M \in \text{Mod-}R$ .

### Corollary of the Baur-Monk Theorem

Let  $M, N \in \text{Mod-}R$ . Then

$$M \equiv N \quad \text{if and only if} \quad |\varphi(M)/\psi(M)| = |\varphi(N)/\psi(N)|,$$

when either is finite, for all  $\varphi \geq \psi \in \text{pp}_R^1$ .

# Pp-types

The **pp-type**,  $\text{pp}^M(\bar{m})$ , of an  $n$ -tuple  $\bar{m}$  in an  $R$ -module  $M$  is the set of all  $\varphi \in \text{pp}_R^n$  such that  $\bar{m} \in \varphi(M)$ .

**Observations:**

- $\text{pp}^M(\bar{m})$  is non-empty.
- If  $\psi \leq \varphi$  in  $\text{pp}_R^n$  and  $\psi \in \text{pp}^M(\bar{m})$  then  $\varphi \in \text{pp}^M(\bar{m})$ .
- If  $\varphi_1, \dots, \varphi_n \in \text{pp}^M(\bar{m})$  then  $\bigwedge_{i=1}^n \varphi_i \in \text{pp}^M(\bar{m})$ .

Therefore  $\text{pp}^M(\bar{m})$  is a filter in  $\text{pp}_R^n$ .

# Free realisations

## Definition (Prest)

Let  $\varphi$  be a pp- $n$ -formula. A **free realisation** of  $\varphi$  is a pair  $(M, \bar{m})$  where  $M$  is a finitely presented  $R$ -module and  $\bar{m}$  is an  $n$ -tuple of elements of  $M$  such that for all  $\sigma \in \text{pp}_R^n$ ,  $\sigma \geq \varphi$  if and only if  $\bar{m} \in \sigma(M)$ .

## Theorem

- (i) *Let  $M$  be a finitely presented  $R$ -module and  $\bar{m}$  an  $n$ -tuple from  $M$ . Then  $\text{pp}^M(\bar{m})$  is generated as a filter by some  $\varphi \in \text{pp}_R^n$  i.e.  $(M, \bar{m})$  is a free realisation of  $\varphi$ .*
- (ii) *Let  $\varphi \in \text{pp}_R^n$ . There exists a finitely presented  $R$ -module and an  $n$ -tuple  $\bar{m}$  such that  $(M, \bar{m})$  is a free realisation of  $\varphi$ .*

## The Compactness Theorem

Let  $R$  be a ring and let  $\mathcal{L}_R := (0, +, (\cdot r)_{r \in R})$  be the language of  $R$ -modules.

- (ii) Let  $\Sigma$  be a set of  $\mathcal{L}_R$ -formulas with free variables  $(v_1, v_2, \dots, v_n)$ .  
If for every finite subset  $\Sigma'$  of  $\Sigma$ , there exists an  $R$ -modules  $M$  and a tuple of elements  $\bar{m} \in M$  such that  $\theta(\bar{m})$  holds in  $M$  for all  $\theta \in \Sigma'$   
then there exists an  $R$ -module  $M$  and a tuple of elements  $\bar{m} \in M$  such that  $\theta(\bar{m})$  holds in  $M$  for all  $\theta \in \Sigma$ .

# Application of the Compactness Theorem

## Proposition

*Any filter  $p$  in  $\text{pp}_R^n$  is the  $pp$ -type of an element of some  $R$ -module.*

## Proof ( $n=1$ ).

Let  $\Sigma := \{\varphi(x) \mid \varphi \in \text{pp}_R^1 \text{ with } \varphi \in p\} \cup \{\neg\psi(x) \mid \psi \in \text{pp}_R^1 \text{ with } \psi \notin p\}$ .

By the Compactness Theorem, it is enough to show that for all

$$\varphi_1, \dots, \varphi_k \in p \text{ and } \psi_1, \dots, \psi_l \notin p,$$

there exists  $M \in \text{Mod-}R$  and  $m \in M$  such that

$$m \in \varphi_i(M) \text{ for all } 1 \leq i \leq k \text{ and } m \notin \psi_i(M) \text{ for all } 1 \leq i \leq l.$$

Since  $p$  is a filter,  $\varphi := \varphi_1 \wedge \dots \wedge \varphi_k \in p$  and  $\psi_i \not\leq \varphi$  for  $1 \leq i \leq l$ .

Let  $(M, m)$  be a free realisation of  $\varphi$ . Then, for all  $\sigma \in \text{pp}_R^1$ ,

$$m \in \sigma(M) \text{ if and only if } \sigma \geq \varphi.$$

Therefore  $m \in \varphi_i(M)$  for all  $1 \leq i \leq k$  and  $m \notin \psi_i(M)$  for  $1 \leq i \leq l$ .  $\square$

# Modules up to Elementary Equivalence

## Definition

An embedding  $f : M \hookrightarrow N$  is **pure** if for all pp-1-formulae  $\varphi$ ,

$$\varphi(N) \cap f(M) = f(\varphi(M)).$$

An  $R$ -module  $M$  is **pure-injective** if every pure-embedding  $M \hookrightarrow N$  splits.

Write  $\text{pinj}_R$  for the set of indecomposable pure-injective  $R$ -modules.

## Theorem (Ziegler)

For every  $R$ -module  $M$ , there exists  $N_i \in \text{pinj}_R$  such that  $M$  is elementary equivalent to  $\bigoplus_{i \in I} N_i$ .

## Theorem

Let  $N, M \in \text{pinj}_R$ . Then  $N \equiv M$  if and only if for all  $\varphi \geq \psi \in \text{pp}_R^1$

$$|\varphi(N)/\psi(N)| > 1 \Leftrightarrow |\varphi(M)/\psi(M)| > 1.$$

# The Ziegler Spectrum

The **Ziegler spectrum**,  $Zg_R$ , of  $R$  is the topological space with set of points  $\text{pinj}_R$  and basis of open sets

$$(\varphi/\psi) := \{N \in \text{pinj}_R \mid |\varphi(N)/\psi(N)| > 1\}$$

where  $\varphi \geq \psi \in \text{pp}_R^1$ .

Definable subcategories of  $\text{Mod-}R$  are in bijective correspondence with the closed subsets of the Ziegler spectrum via the map

$$\mathcal{D} \mapsto \mathcal{D} \cap Zg_R.$$

## Properties of $Zg_R$

- $N, M \in \text{pinj}_R$  are topologically indistinguishable if and only if they are elementary equivalent.
- The sets  $(\varphi/\psi)$  are compact. In particular,  $Zg_R = (x = x/x = 0)$  is compact.
- $Zg_R$  is often not  $T_0$  and very rarely has the property that the intersection of two compact open sets is compact.



## An example

The indecomposable pure-injective modules over  $\mathbb{Z}$  are

- the finite modules  $\mathbb{Z}/p^n\mathbb{Z}$  for  $p$  prime and  $n \in \mathbb{N}$ ,
- the  $p$ -adic integers  $\overline{\mathbb{Z}}_{(p)}$  for  $p$  prime,
- the  $p$ -Prüfer group  $\mathbb{Z}_{p^\infty}$  for  $p$  prime, and
- $\mathbb{Q}$ , the field of fractions of  $\mathbb{Z}$ .

A subset  $\mathcal{C}$  of  $\text{Zg}_{\mathbb{Z}}$  is closed if and only if the following conditions hold:

- If  $\mathcal{C}$  contains infinitely many finite modules then  $\mathcal{C}$  contains  $\mathbb{Q}$ .
- If  $\mathcal{C}$  contains infinitely many  $\mathbb{Z}/p^n\mathbb{Z}$  for fixed prime  $p$  then  $\mathcal{C}$  contains  $\overline{\mathbb{Z}}_{(p)}$  and  $\mathbb{Z}_{p^\infty}$ .
- If  $\mathcal{C}$  contains  $\overline{\mathbb{Z}}_{(p)}$  or  $\mathbb{Z}_{p^\infty}$  then  $\mathcal{C}$  contains  $\mathbb{Q}$ .

# Soberness and the Baire property

## Definition

We say a topological space has the **Baire property** if every countable intersection of open and dense subsets is dense.

## Theorem (Herzog - reinterpreted)

*Every closed subset of  $Zg_R$  has the Baire property.*

## Definition

Let  $\mathcal{T}$  be a topological space.

1. A subset  $S$  of  $\mathcal{T}$  is irreducible if whenever  $S \subseteq C_1 \cup C_2$  where  $C_1$  and  $C_2$  are closed subsets then  $S \subseteq C_1$  or  $S \subseteq C_2$ .
2.  $\mathcal{T}$  is sober if every non-empty irreducible closed set is the closure of a point.

## Corollary (Herzog)

*If  $Zg_R$  has a countable basis then  $Zg_R$  is sober. In particular, if  $R$  is countable then  $Zg_R$  is sober.*

## Remark

*If  $\mathcal{T}$  is a topological space with a countable basis of open sets such that every closed subset of  $\mathcal{T}$  has the Baire property then  $\mathcal{T}$  is sober.*

## Proof.

A topological space  $\mathcal{V}$  is the closure of a point  $x$  if and only if  $x$  is a member of every non-empty open subset of  $\mathcal{V}$ .

Let  $\mathcal{V}$  be an irreducible topological space. Every non-empty open subset  $\mathcal{U}$  of  $\mathcal{V}$  is dense because if  $\mathcal{U} \cap \mathcal{U}' = \emptyset$  for  $\mathcal{U}' \subseteq \mathcal{V}$  open then

$$(\mathcal{V} \setminus \mathcal{U}) \cup (\mathcal{V} \setminus \mathcal{U}') = \mathcal{V}.$$

Therefore if  $\mathcal{V}$  has the Baire property and a countable basis of open sets then  $\mathcal{V}$  is the closure of a point. □

# Soberness and Duality

We write  ${}_R\text{pp}^n$  for the lattice of left pp- $n$ -formulae and  ${}_R\text{Zg}$  for the left Ziegler spectrum of  $R$ .

## Duality for pp-formulae (Prest)

For each  $n \in \mathbb{N}$ , there is an order anti-isomorphism

$$D : \text{pp}_R^n \rightarrow {}_R\text{pp}^n.$$

For any topological space  $\mathcal{T}$ , the open subsets of  $\mathcal{T}$  are a (complete) lattice, denoted  $\mathcal{O}(\mathcal{T})$ , under inclusion.

## Duality for Ziegler Spectra (Herzog)

The map on basic open subsets of  $\text{Zg}_R$ , defined by

$$(\varphi/\psi) \mapsto (D\psi/D\varphi),$$

induces a lattice isomorphism from  $\mathcal{O}(\text{Zg}_R)$  to  $\mathcal{O}({}_R\text{Zg})$ .

This implies that if  $\text{Zg}_R$  and  ${}_R\text{Zg}$  are sober there is a homeomorphism

$$\text{Zg}_R/T_0 \rightarrow {}_R\text{Zg}/T_0.$$

# Is the Ziegler Spectrum always sober?

We don't know.

## Remark

*If an irreducible closed set  $\mathcal{C}$  contains a point  $x$  which is isolated in  $\mathcal{C}$  then  $\mathcal{C}$  is equal to the closure of  $x$ . Hence, if a topological space has Cantor-Bendixson rank then it is sober.*

## Theorem (Gregory-Puninski)

*If  $R$  is a Prüfer ring then  $Zg_R$  is sober. If  $R$  is a (uni)serial ring then  $Zg_R$  is sober.*

## Remark/Theorem

If  $\mathcal{A}$  is a tubular algebra then  $Zg_{\mathcal{A}}$  is sober.

A ring  $R$  is von Neumann regular if for all  $a \in R$  there exists  $x \in R$  such that  $a = axa$ .

## Lemma

*If  $R$  is a von Neumann regular ring then  $Zg_R$  is sober if and only if for all prime ideals  $P$ , there exists an irreducible right ideal  $I$  such that  $P$  is the largest two-sided ideal contained in  $I$ .*

-Thank you-

## An example

Let  $\aleph_1$  be the set of all countable ordinals. The subsets  $(\alpha, \aleph_1) \subseteq \aleph_1$  are the open sets of a topology. The closed sets are  $[0, \alpha]$  for  $\alpha \in \aleph_1$  and  $\aleph_1$ . They are all irreducible.

The sets  $(\alpha, \aleph_1)$  are compact because

$$(\alpha, \aleph_1) \subseteq \bigcup_{\beta \in I} (\beta, \aleph_1)$$

if and only if  $\beta \leq \alpha$  if and only if  $(\alpha, \aleph_1) \subseteq (\beta, \aleph_1)$  for some  $\beta \in I$ .

Every closed subset of  $\aleph_1$  has the Baire property but  $\aleph_1$  is not the closure of a point because  $\bigcap_{\alpha \in \aleph_1} (\alpha, \aleph_1) = \emptyset$ .

# Fun with Ziegler spectra of Artin algebras

Let  $\mathcal{A}$  be an Artin algebra.

The indecomposable finite length modules are exactly the isolated points in  $\text{Zg}_{\mathcal{A}}$ . The set of indecomposable finite length modules is dense in  $\text{Zg}_{\mathcal{A}}$ .

## Theorem

*If  $\mathcal{A}$  is not finite representation type then there is a infinite length indecomposable pure-injective  $\mathcal{A}$ -module.*

## Theorem (Herzog)

*If there are infinitely many indecomposable modules of endolength  $n$ , then there is an infinite length indecomposable module of endolength  $\leq n$ .*

## Proof.

The set of indecomposable modules of endolength  $\leq n$  is a closed (and hence compact) subset of  $\text{Zg}_{\mathcal{A}}$ . Since this set contains infinitely many points, it must contain a non-isolated point. □